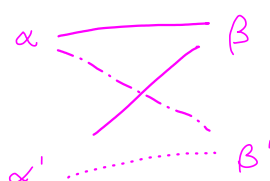


Recall from last lecture:

Sets	Items	Problem	Notation
Matchings	Edges	Finding maximum	$\alpha'(G)$
Edge covers	Edges	Finding minimum	$\beta'(G)$
Vertex covers	Vertices	Finding minimum	$\beta(G)$
Independent set	Vertices	Finding maximum	$\alpha(G)$



We still need to find some relations between α' and β' , and between α and β' (which will be possible under some conditions only).

Theorem (Gallai, 1959, $\alpha' - \beta'$)

If G is a graph without isolated vertices and n vertices in total, then $\alpha'(G) + \beta'(G) = n$. (max. matching + min edge cover)

Proof

2 steps:

1. From a maximum matching (of size $\alpha'(G)$), construct a minimum cover of size $n - \alpha'(G)$. That implies that $n - \alpha'(G) \geq \beta'(G)$.
2. From a minimum cover (of size $\beta'(G)$), construct a matching of size $n - \beta'(G)$. That implies that $\alpha'(G) \geq n - \beta'(G)$.

These two steps prove that $\alpha'(G) + \beta'(G) = n$.

Step 1: Let M be a maximum matching, and let S be a set of vertices. At first, $S = M$. S will be, at the end of the process, an edge cover. For every unsaturated vertex, add to S an edge incident to it.

Claim: S is a cover of size $n - |M|$.

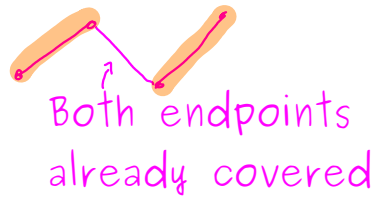
– S is a cover, because edges of M cover saturated edges, and we

added edges to cover unsaturated edges.

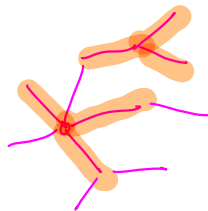
- S has size $n - |M|$: we know S has size $|M| + \text{"number of unsaturated vertices"}$. However, the number of unsaturated vertices is $n - 2|M|$, because every edge of the matching saturates 2 vertices, with no edges saturating the same vertex. So $|S| = |M| + n - 2|M| = n - |M|$.

Let S be a minimum cover. We want to construct a matching of size $n - |S|$.

Observation: In a minimum cover, there is no path of length 3.

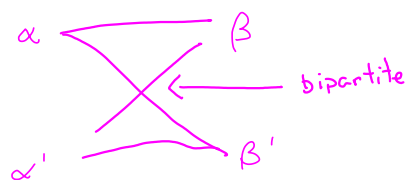


Since there is no path of length 3 in a cover, the components of a cover are all stars. Assume there are k components. Then we can create a matching of size k , by taking one edge in every star. To find the value of k , we count the number of centers of the star. For each edge of the cover, there is exactly one non-center (of a star). So the number of stars is $n - \text{"non-centers"} = n - |S|$. So there is a matching of size $n - |S|$.



Corollary (König, 1916, $\alpha = \beta'$)

If G is a bipartite graph with no isolated vertex, then the size of a maximum independent set is the size of a minimum edge cover.



A vertex cut (or separating set) is a subset of vertices S such that $G-S$ has more than one component.

The connectivity of G , $\kappa(G)$, is the minimum size of a separating set, if it exists, or $n-1$.

A graph is k -connected if its connectivity is at least k .

Examples

Disconnected = connectivity 0

Connected = 1-connected

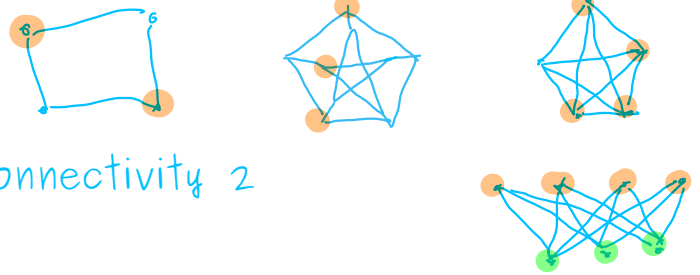
Cycles of length at least 3 have connectivity 2

Petersen graph has connectivity 3.

Complete graph K_n has connectivity $n-1$.

Complete bipartite graph $K_{m,n}$ has connectivity $\min\{m, n\}$.

By convention, we say the graph with one vertex has connectivity 0.



Proposition

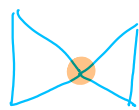
The connectivity of a connected graph is at most its minimum degree.

Proof

One can isolate a single vertex removing all the vertices around it.

Remark

The connectivity of a connected graph is not at least its minimum degree.



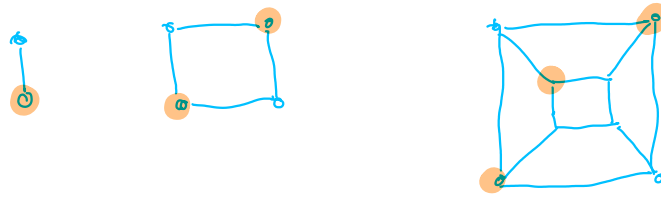
Minimum degree 2, but there is a cut-vertex \Rightarrow connectivity 1.

Example

The hypercube H_k has connectivity k .

Of course, since it is k -regular, it has connectivity at most k .

We can prove by induction it has connectivity at least k :



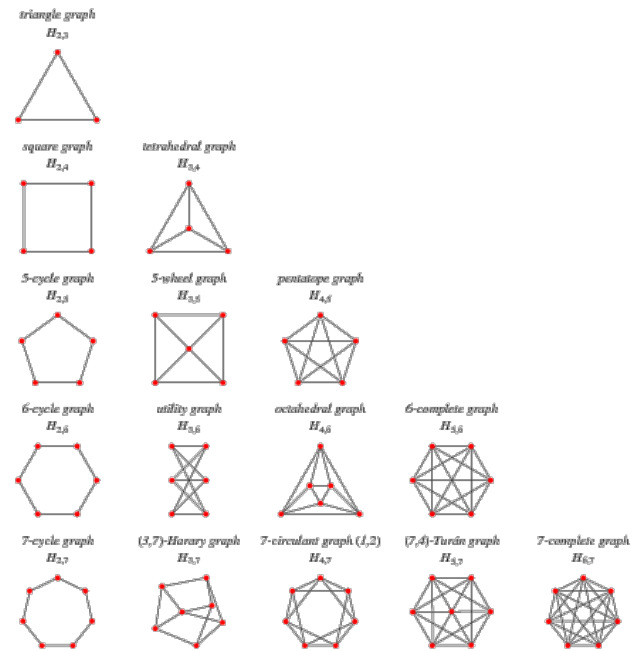
Example: Harary graphs

Harary graphs $H_{n,k}$ are graphs with n vertices and $\lceil \frac{nk}{2} \rceil$ edges, being as regular as possible.

They have connectivity k :

– k is the minimum degree of $H_{n,k}$

Homework: Read the proof in the ^{n,k} text-book that it has connectivity at least k .



Corollary (Harary, 1962)

The minimum number of edges in a k -connected graph with n vertices is $\lceil \frac{nk}{2} \rceil$

Reference: Douglas B. West. Introduction to graph theory, 2nd edition, 2001. Section 3.1 (covers) and section 4.1 (connectivity)