Math 38 - Graph Theory Nadia Lafrenière Chromatic polynomial for chordal graphs 05/27/2020

The goal of today's lecture is to give a shortcut for computing the chromatic polynomial of many graphs.

Geometry

Simplex (triangle, tetrahedron, etc.)





Simplicial complex (gluing simplices together) Every face of a simplicial complex is a simplex.



Graph theory

Clique







Chordal graph, i.e. graph that admit a simplicial elimination ordering.



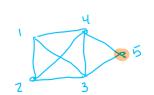


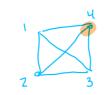
A vertex is simplicial if its neighbors form a clique.



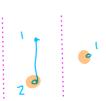
- Simplicial
- Not simplicial

A <u>simplicial elimination ordering</u> is an ordering of the vertices v_n , ..., v_n such that the vertex v_i is simplicial in $G - \{v_n, ..., v_n\}$









The ordering above is a simplicial elimination ordering.

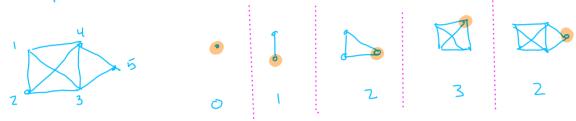
Example

- 2
- For complete graphs, any ordering is a simplicial elimination ordering.
- For trees, an ordering starting at the leaves and going inwards is a simplicial elimination ordering.
- The cycle of length 4 has no simplicial ordering.

Proposition

The chromatic polynomial of a graph with simplicial ordering $v_{\mu}...,v_{\nu}$ is the product of the k-d'(v_{μ})'s, where d'(v_{μ}) is the degree of v_{μ} in the induced subgraph with vertices $\{v_{\mu}...,v_{\mu}\}$.

Example



$$\chi(G;k) = k (k-1)(k-2)^2(k-3)$$

Sketch of proof

When it is added, v_i has $d'(v_i)$ options for coloring it. Doing this process, we count all the options for coloring the graph with k colors.

Chordal graphs

A chord of a cycle C is an edge not in C that whose endpoints are in C



Not a chorc



A cycle is <u>chordless</u> if it has length at least 4 and no chord. A graph is chordal if it is simple and it has no chordless cycle (as induced subgraphs).

Theorem (Dirac, 1961)

A simple graph has a simplicial elimination ordering if and only if it is a chordal graph.

Lemma (Voloshin, 1982, or Farber-Jamison, 1986) Every chordal graph has a simplicial vertex.

The proof of the lemma is omitted. It is in the textbook as Lemma 5.3.16.

Proof of the theorem

Necessity: If it is not chordal, there is a chordless cycle C. In C, that has length at least 4, none of the vertices is simplicial, and no vertex can be the first one to be picked in C for the ordering. Sufficiency: Using the lemma, we know that every chordal graph G has a simplicial vertex v. Delete v. Then $G-\{v\}$ is chordal, so we can apply

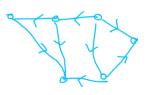
Acyclic orientations

What is the meaning of $\chi(G;-1)$?

the lemma again, creating an ordering.

An <u>orientation</u> of a graph G is a digraph that has G as an underlying graph.





orientation

acyclic orientation

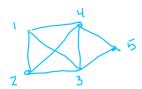
An orientation is acyclic if it has no cycle.

Theorem (Stanley, 1972)

The absolute value of $\chi(G;-1)$ is the number of acyclic orientations of G.

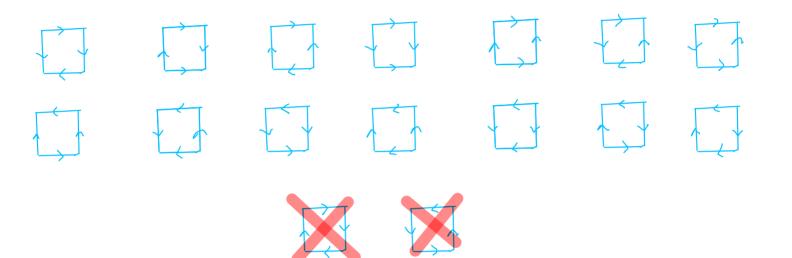
Homework Read proof of Theorem 5.3.27 in the textbook.

Examples



has chromatic polynomial $\chi(G;k)=k(k-1)(k-2)^2(k-3)$, so if has 72 acyclic orientations.

 C_{4} has chromatic polynomial $\chi(C_{4};k)=k(k-1)(k^2-3k+3)$, and $\chi(C_{4};-1)$ is 14.



Reference: Douglas B. West. Introduction to graph theory, 2nd edition, 2001. Section 5.3.