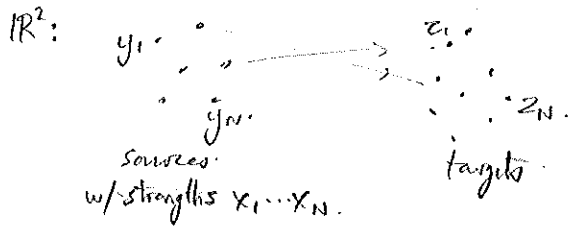


last time empirically observed low-rank property of off-diag blocks, but how exploit this? (without doing an SVD, which is  $O(N^3)$ )

last time: want eval potential  $u(z) = \sum_{j=1}^N x_j \ln \frac{1}{|z - y_j|}$  at  $z = z_i, i = 1 \dots N$ .



change strengths of sources  
 • since  $N$  sources,  $N$  targets,  $N^2$  interactions (have to eval  $\ln$  dist  $N^2$  times, naively)

• if  $u_i := u(z_i)$  then  $\vec{u} = \tilde{A} \vec{x}$ , where  $\tilde{A}$  is some  $N \times N$  off-diagonal block of e.g.  $N$ -point matrix.

Note:  $u(z)$  harmonic for  $z \neq y_j, j = 1 \dots N$ .

Thm (multipole expansion): Let  $B$  be a disc containing all  $y_j$ , centered at  $O$ , radius  $R$ ,

then for  $r > R$ ,  $u(r, \theta) = c_0 \ln \frac{1}{r} + \sum_{n=1}^{\infty} (a_n \cos n\theta + b_n \sin n\theta) r^{-n}$   
 $z$  in polar  $\nearrow$   
 sums abs. convergent in  $r > R$ .  
 Note: setting  $c_0 = 0$ , true for any func  $u$  harmonic in  $r \geq R$  w/  $u = o(1)$  as  $r \rightarrow \infty$ .

Or considering  $\mathbb{R}^2 \simeq \mathbb{C}$ ,  $u(z) = \text{Re} \left[ c_0 \ln \frac{1}{z} + \sum_{n=1}^{\infty} c_n z^{-n} \right]$   
 Laurent expansion ("Taylor exp. about  $\infty$ ")

Say truncate to  $p$  terms, how bad is error?

Consider single unit charge @  $y$ :  $u(z) = \ln \frac{1}{z-y}$  let's work w/ complex-valued potential, take  $\text{Re}$  at end.  
 $= \ln \frac{1}{z} - \ln(1 - y/z)$   
 $= \ln \frac{1}{z} + yz^{-2} + \frac{y^2}{2} z^{-3} + \frac{y^3}{3} z^{-4} + \dots$  abs. conv. for  $|z| > |y|$   
 is multipole expansion, proves above thm. Taylor  $-\ln(1-x) = x + x^2/2 + x^3/3 + \dots$  when  $|x| < 1$ .

Pointwise truncation error  $e_p(z) := \underbrace{\ln \frac{1}{z} + \sum_{n=1}^{p-1} \frac{y^n}{n} z^{-n}}_{p\text{-term approx.}} - \underbrace{\ln \frac{1}{z-y}}_{\text{true}} = \sum_{n=p}^{\infty} \frac{y^n}{n} z^{-n}$  tail of sum.

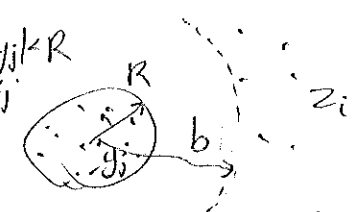
so  $|e_p(z)| \leq \sum_{n=p}^{\infty} \frac{1}{n} \left| \frac{y}{z} \right|^n = \left| \frac{y}{z} \right|^p \sum_{n=0}^{\infty} \frac{1}{n+p} \left| \frac{y}{z} \right|^n \leq \left| \frac{y}{z} \right|^p \left( \frac{1}{p} + \ln \frac{1}{1 - |y/z|} \right)$   
 shift sum.  $\underbrace{\quad}_{\leq 1/n \text{ for } n > 0}$   
 so  $|e_p(z)| = O\left(\left| \frac{y}{z} \right|^p\right)$  as  $p \rightarrow \infty$  exponential conv. in  $p$ !  $O(1)$  as  $p \rightarrow \infty$ , for fixed  $y, z \in \mathbb{C}$ .

Use this bound for each charge in disc  $B$ , get:

Thm (multipole error): potential due to  $N$  charges  $x_j$ , locations  $y_j$ , in disc rad.  $R$  can be rep. by  $p^{\text{th}}$ -order multipole expansion in  $|z| > b > R$  w/ pointwise error  $\leq C \left( \sum_{j=1}^N |x_j| \right) \cdot \left( \frac{R}{b} \right)^p$

Use to apply off-diagonal blocks  $\tilde{A}$ :  $|z_i| > b$ ,  $|y_j| < R$

Say  $\{z_i\}_{i=1}^N$  'well-separated' from  $\{y_j\}_{j=1}^N$ , i.e.  $\frac{b}{R}$  significantly larger than 1, eg. 2.



Recipe: (HW7!)  
 (i) decide  $p$  based on desired accuracy:  $(\frac{R}{b})^p$  (tot. change)  $\approx \epsilon$ .  $\leftarrow$  typ.  $p \sim 20$  to  $40$ .  
 (ii) compute multipole coeffs due to sources:  

$$C_0 = \sum_{j=1}^N x_j, \quad C_n = \sum_{j=1}^N \frac{y_j^n}{n} x_j, \quad n=1, \dots, p-1$$
  
 (iii) evaluate multipole expansion at targets:  

$$U(z_i) \approx U^{(p)}(z_i) = C_0 \ln \frac{1}{z_i} + \sum_{n=1}^{p-1} C_n z_i^{-n} \quad i=1, \dots, N.$$

Complexity:  $N \times \tilde{A} \approx N \times \begin{matrix} p \\ \times \\ p \end{matrix} \times N$   
 compare original  $O(N^2)$ ! eg.  $N=10^6$ ,  $p=10^{1.5}$  ( $\epsilon=10^{-9}$ ) then speedup is  $10^{4.5} \approx 30000$ ! That's a good algorithm!

Unfortunately, applying whole of  $A$ , interaction matrix between  $\{y_j\}_{j=1}^N \leftarrow$  i.e. sources = targets is trickier since not all clumps of points well-separated.

Say all  $y_j$ 's lie in rectangle, roughly uniformly distributed:

Cover by  $M$  square boxes: box  $B_i$  all these boxes are well-separated from  $B_i$   
 if uniform,  $\approx \frac{N}{M}$  charges per box.  
 $b = \frac{3}{2}L$ ,  $R = \frac{L}{\sqrt{2}}$  }  $\frac{b}{R} = \frac{3\sqrt{2}}{2} \approx 2.1$  Controls convergence rate.

Effort to get multipole coeffs about each box due to charges  $n$  it =  $pN$   $\leftarrow$  tot charges terms each charge affects.  
 $z - z_0$  replaces  $z$  in expansion, where  $z_0 =$  box center.

Then  $U_i = \sum_{j=1}^N A_{ij} x_j = \sum_{\substack{j \in B \\ \text{or touching box}}} A_{ij} x_j + \sum_{\substack{j \text{ in box for} \\ \text{which } B \text{ is well-sep.}}} A_{ij} x_j$   
 do direct sum, effort  $\approx q \frac{N}{M}$   
 approx by sum of multipole expansions from each of  $M-q$  other boxes, effort  $\approx p(M-q) = O(pM)$ .

Total effort ( $i=1, \dots, N$ )  
 $= \underbrace{pN}_{\text{get coeffs (cheap)}} + \underbrace{q \frac{N^2}{M}}_{\text{local, direct}} + \underbrace{pMN}_{\text{distant.}}$   
 $\leftarrow$  how choose  $M$  to scale with  $N$  so best?

optimal here is  $O(pN^{3/2})$   
 Fixing  $p$ , this is  $O(\frac{N^{1/2}}{p})$  times faster than naive, eg.  $N=10^6$ ,  $p=10^{1.5}$ , is  $10^{1.5} \approx 30 \times$  faster.  
 Ok, but not as great as before.