Abstract. Let $\sigma(n)$ denote the sum of all of the positive divisors of $n$, and let $s(n) = \sigma(n) - n$ denote the sum of the proper divisors of $n$. The functions $\sigma(\cdot)$ and $s(\cdot)$ were favorite subjects of investigation by the late Paul Erdős. Here we revisit three themes from Erdős’s work on these functions. First, we improve the upper and lower bounds for the counting function of numbers $n$ with $n$ deficient but $s(n)$ abundant, or vice versa. Second, we describe a heuristic argument suggesting the precise asymptotic density of $n$ not in the range of the function $s(\cdot)$; these are the so-called nonaliquot numbers. Finally, we prove new results on the distribution of friendly $k$-sets, where a friendly $k$-set is a collection of $k$ distinct integers all of which share the same value of $\sigma(n)/n$.

1. Introduction

Let $\sigma(n)$ be the sum of the natural number divisors of $n$, so that $\sigma(n) = \sum_{d|n} d$. Let $s(n)$ be the sum of only the proper divisors of $n$, so that $s(n) = \sigma(n) - n$. Interest in these functions dates back to the ancient Greeks, who classified numbers as deficient, perfect, or abundant according to whether $s(n) < n$, $s(n) = n$, or $s(n) > n$, respectively. For example, 10 is deficient, 12 is abundant, and 6 is perfect. The ancient Greeks also found and remarked on 2-cycles, where $s(n) = m$ and $s(m) = n$, calling such a pair amicable; for example, 220 and 284.

A conjecture of Catalan in 1888 [8], as amended by Dickson in 1913 [12], claims that for every positive integer $n$, the aliquot sequence $n, s(n), s(s(n)), \ldots$ either terminates at 0 or enters some cycle, such as a perfect number, an amicable pair, or some longer cycle. Despite much numerical work, there are still 12 numbers $n < 1000$ for which the conjecture is in doubt, the first being 276. This numerical work uncovered long stretches where some aliquot sequences grow monotonically and exponentially, and other stretches where they decay in like manner. This behavior was taken into account in the Guy–Selfridge [25] counter-conjecture that for asymptotically all even numbers $n$, the aliquot sequence starting with $n$ is unbounded.

It was proved by Lenstra [28] that for each $k$ there is an aliquot sequence that strictly increases for $k$ terms. Erdős [20] showed that this is common, and in fact, if $n$ is abundant, then but for a set of numbers $n$ of asymptotic density 0, the aliquot sequence starting with $n$ strictly increases for $k$ terms. In the same paper, Erdős claimed the analogous result for deficient numbers (with “increases” replaced with “decreases”), but this claim was retracted in [22]. In that paper, the claim was shown for $k = 2$. It is conjectured that the full claim holds, but this has never been shown.

These results assert that if $n$ is abundant, it is likely that $s(n)$ is also abundant, and similarly for deficient. Say $n$ is an aliquot reversal if this fails. More precisely, say $n$ is an up-down reversal if $n$ is nondeficient yet $s(n)$ is deficient, and say $n$ is a down-up reversal if $n$ is deficient and $s(n)$ is nondeficient.

Though the terminology is new to this paper, aliquot reversals have been studied for some time because of their connection with amicable pairs; in fact the lesser member of an amicable pair is an up-down reversal, while the larger member is a down-up reversal.
Erdős was the first to prove, in [15], that the set of numbers involved in an amicable pair has vanishing asymptotic density, and he did this by showing that there are few up-down reversals. The count of up-down reversals was improved later by Rieger [39], by Erdős–Rieger [21], and by the second author (see [35]). In that last paper it was shown that the number of \( n \) giving an up-down reversal is at most

\[
x/ \exp((c + o(1))(\log_3 x \log_4 x)^{1/2}), \text{ as } x \to \infty,
\]

for a certain absolute constant \( c > 0 \). Here, the subscripts indicate iteration of the natural log function. The proof uses results of Erdős [13] for the count of primitive nondeficient numbers. Incorporating an improvement on these results due to Avidon [3], this argument gives \( c = 1 \). In [32, Corollary 1.5], the same upper estimate (1) was shown for the number of down-up reversals in \([1, x]\), but only with the smaller constant \( c = \frac{1}{10} \).

Our first theorem is an improvement on these upper bounds.

**Theorem 1.1.** Both the count of down-up reversals in \([1, x]\) and the count of up-down reversals in \([1, x]\) are bounded above by (1) with \( c = \sqrt{3} \).

The proof uses recent results from [2] on the number of solutions \( n \) to congruences of the form \( \sigma(n) \equiv a \pmod{n} \).

The problem of obtaining lower bounds on the number of aliquot reversals does not seem to have been previously considered. Note that the upper bounds of the last paragraph just barely give density 0. In the following results we show there is a reason for this: reversals are fairly common. We prove the following.

**Theorem 1.2.** The number of up-down reversals in \([1, x]\) is \( \gg x/(\log_2 x)(\log_3 x)^3 \).

**Theorem 1.3.** The number of down-up reversals in \([1, x]\) is \( \gg x/(\log_2 x)(\log_3 x)^2 \).

One might wonder if these results lend support to the Catalan–Dickson conjecture or to the Guy–Selfridge counter conjecture. It is helpful to think in terms of parity. In fact, \( n \) and \( s(n) \) have opposite parity precisely when \( n \) is a square or a double of a square, and such \( n \) comprise a very sparse set. In addition, it follows from the proof in [20] that for each fixed \( k \), almost all \( n \) have the same parity as \( s_k(n) \) (the \( k \)-th iterate of \( s \) at \( n \)). Since an aliquot sequence terminates only if it hits a prime, a perhaps reasonable conjecture is that aliquot sequences beginning with even numbers usually either are unbounded or enter some cycle of even numbers. Numbers involved in such cycles have density 0 (see [27]). Numerical evidence indicates that most numbers involved in these cycles are amicable (i.e., they are in a 2-cycle), and we know that amicable numbers are quite sparsely distributed (see [37]).

Thus, despite the fact that our results show non-monotonicity is fairly common, the Guy–Selfridge point of view is believable. And some cases of non-monotonicity are quite rare. For example, take numbers \( n \equiv 6 \pmod{12} \), \( n > 6 \). They are abundant, but the number of them to \( x \) with \( s(n) \) deficient is \( O(x^{1/2}) \). Further, in a recent preprint of Bosma [6], it is shown that about \( 1/3 \) of the even numbers to 1 million have their aliquot sequence still open beyond \( 10^{100} \). This evidence seems to put a chill on the Catalan–Dickson conjecture!

On the other hand, it is hard to align this numerical evidence with the recent paper of Bosma–Kane [7] where it is shown that the geometric mean of \( s(2n)/2n \) for \( n \leq x \) tends to a constant smaller than 1 as \( x \to \infty \).

We remark that the lower bounds in Theorems 1.2, 1.3 are established for even numbers. It would be of interest to see what kind of lower bounds can be established for odd numbers.

We turn now to our second topic. Positive integers not in the range of \( s(\cdot) \) are said to be *nonaliquot*, though the more colorful term *untouchable* (see Alaen [1]) has also been used. The concept dates back to at least ca. 1000 CE [40]. A useful initial observation is that if \( p \) and \( q \) are distinct primes, then \( s(pq) = p + q + 1 \). A slightly stronger form of Goldbach’s conjecture is that every even integer \( n \geq 8 \) is a sum of two distinct primes. If so, then every odd \( n \geq 9 \) belongs to the range of \( s(\cdot) \); noting that \( s(2) = 1, s(4) = 3, \) and \( s(8) = 7 \), we would have that 5 is the only odd nonaliquot number.
While the conjecture of the last paragraph remains out of reach, it follows from (independent) work of Chudakov, van der Corput, and Estermann that almost all even numbers can be written as a sum of two distinct primes. (See [45, Chapter 3] for a modern treatment.) Hence, the set of odd nonaliquot numbers has asymptotic density zero.

What about even numbers? In [18], Erdős proved that a positive proportion of even numbers are nonaliquot. By an elaboration of Erdős’s methods, Chen and Zhao [9] have shown that the nonaliquot even integers make up a set of lower density greater than 0. (This improves earlier results reported in [5] and [43, Corollary 9.2, p. 64].) Erdős was unable to decide whether or not a positive proportion of even numbers do belong to the image of \( s \), but this has been resolved affirmatively in very recent work of Luca and the second author [29].

In §3, refining some thoughts in [43], we propose a probabilistic model for \( s(x) \) that suggests a precise value for the density of nonaliquot numbers.

**Conjecture 1.4.** The set of nonaliquot numbers has asymptotic density \( \Delta \), where

\[
\Delta = \lim_{y \to \infty} \frac{1}{\log y} \sum_{2 \leq a \leq y} \frac{1}{a} e^{-a/s(a)}.
\]

At \( y = 2 \cdot 10^{10} \), the expression inside the limit has value 0.171822. Table 1 collects counts of nonaliquot numbers up to \( 10^{10} \); these counts are consistent with a limiting density of \( \approx 0.17 \).

Finally, we consider the distribution of the rational numbers \( \frac{\sigma(n)}{n} \). We are interested here in statistics for

\[
F(n; x) = \# \left\{ m \leq x : m \neq n, \frac{\sigma(m)}{m} = \frac{\sigma(n)}{n} \right\},
\]

thought of as a random variable on the set of natural numbers \( n \leq x \), equipped with the uniform measure. A theorem of Wirsing [46] asserts that an equation of the form \( \frac{\sigma(n)}{n} = \alpha \) has at most \( x^{o(1)} \) solutions \( n \leq x \), as \( x \to \infty \), uniformly in \( \alpha \). Thus, we have a pointwise bound \( F(n; x) = x^{o(1)} \), as \( x \to \infty \). The average behavior of \( F(n; x) \) was determined by Erdős [17, Theorem 2]: for a certain constant \( c > 0 \),

\[
(2) \quad \sum_{n \leq x} F(n; x) \sim cx, \quad \text{as} \quad x \to \infty.
\]

Note that \( \sum_{n \leq x} F(n; x) \) counts the number of ordered pairs of distinct integers \( (n, m) \in [1, x]^2 \) with \( \frac{\sigma(n)}{n} = \frac{\sigma(m)}{m} \), and this is the way Erdős states his result.

<table>
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<th>( x )</th>
<th>( U(x) )</th>
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<th>( x )</th>
<th>( U(x) )</th>
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**Table 1.** Counts \( U(x) \) of nonaliquots to various heights \( x \) from \( 10^8 \) through \( 10^{10} \), along with the corresponding densities \( D(x) \).
Call a pair (unordered) of distinct integers $n$ and $m$ a friendly pair if $\frac{\sigma(n)}{n} = \frac{\sigma(m)}{m}$. Evidently, if $u$ is coprime to $nm$, then $un, um$ is also a friendly pair. Call the friendly pair $n,m$ primitive friendly if $n$ and $m$ have no nontrivial common unitary divisor. (We say $d$ is a unitary divisor of $n$, and write $d \mid n$, if $d \mid n$ and $\gcd(d,n/d) = 1$.) The key ingredient in Erdős’s proof of (2) is the convergence of $\sum \frac{1}{m}$, where $m$ runs over the larger members of the primitive friendly pairs (with multiplicities). We prove a substantially stronger result on the distribution of such pairs:

**Theorem 1.5.** The number of primitive friendly pairs contained in $[1,x]$ is at most $x^{\frac{1}{2}+o(1)}$, as $x \to \infty$.

(The convergence of $\sum \frac{1}{m}$ follows immediately by partial summation.) The exponent $\frac{1}{2}$ appears to be the limit of our method, but the numerical data (see Table 2) suggest that this is not sharp. Perhaps the true count of primitive friendly pairs in $[1,x]$ is $x^{o(1)}$.

By studying primitive friendly $k$-sets for integers $k \geq 2$, we are able to show that $F(n;x)$ has a limiting distribution. Set $N_k(x) = \#\{n \leq x : F(n;x) \geq k\}$.

**Theorem 1.6.** For each fixed nonnegative integer $k$, we have $N_k(x) = (\alpha_k + o(1))x$, as $x \to \infty$, for some constant $\alpha_k$. Moreover, $\alpha_k \to 0$ as $k \to \infty$.

Clearly, $\alpha_0 = 1$. Though the constants $\alpha_k$ for $k \geq 1$ are effectively computable (via the proof of Theorem 1.6), we have not found good rigorous numerical estimates for them. However, explicit computations of $N_k(x)$ suggest that $\alpha_1 = 0.0347\ldots$, $\alpha_2 = 0.0028\ldots$, $\alpha_3 = 0.00085\ldots$, and $\alpha_4 = 0.000084\ldots$; see Table 3.

**Notation.** We use the Bachmann–Landau \(\ll\); \(O\) and \(o\)-notations, with their usual meanings, as well as the associated Vinogradov \(\ll\) and \(\gg\) notations. $P(n)$ denotes the largest prime factor of $n$, with the convention that $P(1) = 1$, and $\rad(n)$ denotes the largest squarefree divisor of $n$. We reserve the letter $p$, with or without subscripts or dashes, for primes. We write $v_p(n)$ for the exponent on the highest power of $p$ dividing $n$. The notation $\gcd(m,n)$ stands for the greatest common unitary divisor of $m$ and $n$. We use $\varphi$ for Euler’s totient function, $\omega$ for the number of prime divisors, $\Omega$ for the number of prime factors counted with multiplicity, and $\tau$ for the total number of divisors.

2. **Aliquot reversals**

2.1. **Upper bounds**
2.1.1. A structural lemma on primitive nondeficient numbers. The proof of Theorem 1.1 runs along similar lines as the proof of the main theorem of [35]. In particular, the primitive nondeficient numbers play a critical role. (These are numbers $n$ with $\sigma(n)/n \geq 2$ and for each proper divisor $d$ of $n$, $\sigma(d)/d < 2$.) The main new ingredient is a lemma that is perhaps of independent interest: Almost all primitive nondeficient numbers $n$ have $P(n) \lesssim n^{1/3}$.

We will deduce the lemma from the results of [2] concerning solutions $n$ to the congruence $\sigma(n) \equiv a \pmod{n}$. In that paper, $n$ is called a regular solution if

$$n = pm, \quad \text{where} \quad p \nmid m, \quad m \mid \sigma(m), \quad \text{and} \quad \sigma(m) = a.$$ 

(It is straightforward to check that these numbers $n$ are solutions of the congruence.) In the remaining case, $n$ is called sporadic. Note that regular solutions can exist only when $a > 0$. The following is a slightly less precise form of the main result of [2].

**Proposition 2.1.** The number of sporadic solutions $n \leq x$ to $\sigma(n) \equiv a \pmod{n}$ is at most $x^{1/2+o(1)}$, as $x \to \infty$, uniformly for $|a| \leq x^{1/4}$.

The restriction $|a| \leq x^{1/4}$ is inconvenient for our purposes. We can ameliorate this as follows. Call a sporadic solution $n$ to $\sigma(n) \equiv a \pmod{n}$ convenient if all of the following hold:

(i) $k := \frac{\sigma(n)-a}{n}$ is an integer at least 2,

(ii) if $d$ is a unitary divisor of $n$ with $d < n$, then $\sigma(d)/d < k$,

(iii) there is a prime $q > P(n)$ with $\frac{\sigma(nq)}{nq} > k$.

**Proposition 2.2.** The number of convenient solutions $n \leq x$ to $\sigma(n) \equiv a \pmod{n}$ is at most $x^{1/2+o(1)}$, as $x \to \infty$, uniformly in $a$ with $|a| \leq x$.

**Sketch of the proof.** The proof in [2] goes by reducing to the case when $n$ is convenient (see that paper’s Lemma 5), in order to locate a divisor of $n$ of size roughly $\sqrt{n}$. If we assume that $n$ is convenient to start with, the proofs of Lemmas 6, 7, and the main theorem can be run essentially as before. (If we mimick the proof of that paper’s Theorem 1, we may now discard the condition that $n > 6a^2 \log(6|a|)$.)

We now deduce the key lemma.

**Lemma 2.3.** Fix $\epsilon > 0$. The number of primitive nondeficient $n \leq x$ with $P(n) > n^{1/3+\epsilon}$ is $O(x^{1-\epsilon})$, as $x \to \infty$.

**Proof.** After a dyadic subdivision argument, it is enough to prove the claimed upper bound for the number of primitive nondeficient $n \in (x, 2x]$ with $P(n) > x^{1/3+\epsilon}$. By the result of Wirsing quoted in the introduction, the number of perfect $n \leq 2x$ is $x^{\alpha(1)}$, as $x \to \infty$. Thus, we can assume that $n$ is abundant. We can also assume $P(n)^2 \nmid n$, since otherwise $n$ belongs to a set of size at most $2x \sum_{p \geq x^{1/3+\epsilon}} x^{-2} \ll x^{2/3-\epsilon}$. Write $n = pm$, where $p = P(m)$. Since $n = mp$ is primitive nondeficient,

$$2 > \frac{\sigma(m)}{m} > \frac{2}{\sigma(p)/p} > 2 - 2x^{-1/3-\epsilon}.$$ 

Hence,

$$-2mx^{-1/3-\epsilon} < a := \sigma(m) - 2m < 0$$

and

$$|a| < 2mx^{-1/3-\epsilon} \leq \frac{4}{p} x^{-1/3-\epsilon} = 4x^{2/3-\epsilon}/p.$$ 

The integer $m$ is a sporadic solution to $\sigma(m) \equiv a \pmod{m}$. Moreover, conditions (i)–(iii) in the definition of a convenient solution are easily verified for $m$, with $k = 2$ and $q = p$. 

Fixing \( a \) and \( p \), Proposition 2.2 shows that the number of corresponding values of \( m \leq 2x/p \) is \( O((x/p)^{1/2+\epsilon/2}) \). Summing on \( a \in (-4x^{2/3-\epsilon}/p, 0) \), the number of possibilities for \( m \) given \( p \) is
\[
\ll \epsilon \frac{x^{2/3-\epsilon}}{p} \cdot \left( \frac{x}{p} \right)^{1/2+\epsilon/2} \ll \frac{x^{7/6-\epsilon}}{p^{3/2}}.
\]
Summing on \( p > x^{1/3+\epsilon} \) gives a final upper bound of \( O(\epsilon x^{1-\epsilon}) \), as desired. \( \square \)

**Remark.** A slightly more careful argument improves the upper bound to \( x^{1-3\epsilon+o(1)} \), as \( x \to \infty \). For our purposes, any power-savings upper bound is sufficient.

It would be desirable to prove Lemma 2.3 with a smaller exponent than \( \frac{1}{3} \). If the lemma holds with \( \frac{1}{R} \) (say), then the arguments below show that Theorem 1.1 holds with \( c = \sqrt{R} \).

In particular, if \( R \) can be taken arbitrarily large, we would have a qualitative improvement over any bound of the shape (1). Such an improvement would follow from a conjecture posed in [2]:

**Conjecture 2.4.** Let \( x \geq 3 \), and let \( a \) be an integer with \( |a| \leq x/2 \). Let \( k \) be an integer with \( k \geq 2 \). The number of solutions \( n \leq x \) to \( \sigma(n) = kn + a \) is
\[
\ll (\log x)^C,
\]
where both the implied constant and \( C \) are absolute constants.

**Remark.** This is a corrected form of the conjecture as it appeared in [2]. The original formulation allowed \( k = 0 \) and \( k = 1 \); these cases must be excluded to account for popular values \( a \) of \( \sigma(n) \) or \( s(n) \).

2.1.2. **Proof of Theorem 1.1.** Let
\[
y = y(x) = \frac{\log_2 x}{2 \log_3 x}.
\]
We begin with a lemma.

**Lemma 2.5.** The number of integers \( n \leq x \) such that \( m \mid \sigma(n) \) for some integer \( m \leq y \) is \( O(x/\log_2 x) \).

**Proof.** We use [35, Theorem 2], which asserts that for coprime integers \( \ell, k \) with \( k > 0 \), the number of integers \( n \leq x \) such that no prime \( p \equiv \ell \mod k \) has \( p \mid n \) is \( O(x/(\log x)^{1/\varphi(k)}) \), uniformly. Applied with \( \ell = -1, k = m \) the result implies that the number of \( n \) counted in the lemma is
\[
\ll \sum_{m \leq y} \frac{x}{(\log x)^{1/\varphi(m)}} \ll \sum_{m \leq y} \frac{x}{(\log x)^{1/m}} \ll \sum_{m \leq y} \frac{x}{(\log_2 x)^2} \ll \frac{x}{\log_2 x}. \quad \square
\]

For future reference we state the result of Avidon [3] mentioned in the introduction.

**Lemma 2.6.** The number of primitive nondeficient numbers \( a \leq t \) is at most
\[
t/\exp((1 + o(1))\sqrt{\log t \log_2 t}), \quad t \to \infty.
\]

We shall also use a result of Toulmonde [44, Théorème 1] (see the discussion in §10) on the modulus of continuity of the distribution function for \( \sigma(n)/n \).

**Lemma 2.7.** If \( \alpha \) is in the range of \( \sigma(n)/n \), then uniformly for \( T \geq 3 \), the number of integers \( m \leq x \) with
\[
\alpha - \frac{1}{T} \leq \frac{\sigma(m)}{m} < \alpha
\]
is \( \ll x/\exp(\frac{1}{5}\sqrt{\log T \log \log T}) \).
Proof of Theorem 1.1. Let $\epsilon > 0$ be fixed. Suppose $n \leq x$ is nondeficient with $s(n)$ deficient, and let $a$ be the least nondeficient number that divides $n$. If each prime power dividing $a$ is at most $y$, then Lemma 2.5 implies that but for a negligible number of choices for $n$, we may assume that $a \mid \sigma(n)$. This in turn implies that $a \mid s(n)$, contradicting that $s(n)$ is deficient. Thus, $a$ has a prime power divisor exceeding $y$. If $a > y^{3-\epsilon}$, then Lemma 2.6 implies that the number of $n \leq x$ divisible by such a primitive nondeficient number $a$ is at most $x/\exp((1+o(1))(3-\epsilon)\log_3 x \log x)$ as $x \to \infty$. Hence we may assume that $a \leq y^{3-\epsilon}$ and $a$ is divisible by a prime power $q > y$. If $q$ is a proper power, the number of $n \leq x$ divisible by such a large power is $O(x/\log_2 x)$, so we may assume that $q$ is prime and $q > a^{1/3+\epsilon/9}$. But then, by Lemma 2.3, the number of $n \leq x$ divisible by such a primitive nondeficient $a$ is $O_e(x/y^{\epsilon/9})$. Since $\epsilon$ may be taken arbitrarily close to 0, our result follows.

The argument for down-up reversals is similar, but more difficult. Suppose $n \leq x$ is deficient and $s(n)$ is not deficient, and let $a$ be the least nondeficient number dividing $s(n)$. Again, by Lemma 2.5, if each prime power dividing $a$ is at most $y$, then we may assume that $a \mid \sigma(n)$. This implies that $a \mid n$, contradicting $n$ deficient. Thus, we may assume that $a$ has some prime power divisor greater than $y$.

Write $n = np$ where $p = P(n)$. By standard results on the distribution of integers all of whose prime factors are small (see [10]), we may assume that $p > x^{1/\log_3 x}$ and $p \not\mid m$. We have $s(n) = ps(m) + \sigma(m) \equiv 0 \pmod{a}$. Let $a' = \gcd(a,s(m))$. The congruence implies that $a' \mid \sigma(m)$, so that $a' \mid m$. The congruence also implies that $p$ is in a residue class modulo $a/a'$.

We first assume that

$$a \leq Z := x^{1/2\log_3 x}.$$ 

Given $a, m, a'$, by the Brun–Titchmarsh inequality, the number of choices for $p \leq x/m$ is

$$\ll \frac{x}{m \varphi(a/a') \log(x/m(a/a'))} \ll \frac{x \log_3 x}{m \varphi(a/a') \log x} \ll \frac{x \log_3 x}{m(a/a') \log x},$$

where for the last step we use that $a$ is primitive nondeficient, so that $\varphi(a/a') \gg a/a'$. We have $a' \mid m$, so summing the displayed upper bound over $m$ and then over $a'$, we get that the number of choices is

$$\ll \frac{\tau(a)}{a} x \log_3 x.$$

We can restrict to $a$ which have no proper prime power divisor in $(y, Z]$. Indeed, by an argument analogous to that just seen, the number of $n$ as above for which $s(n)$ is divisible by a given prime power $\ell^v \leq Z$ is $O(xv \log_3 x / \ell^v)$. Summing over prime powers $\ell^v > y$, we find that only a negligible set of $n$ arise in this manner.

We now wish to find a good upper bound for the sum of $\tau(a)/a$ as $a$ runs over primitive nondeficient numbers divisible by a prime power greater than $y$, and satisfying the restriction of the last paragraph. We first assume that $a \leq y^{3-\epsilon}$. Then $P(a) > a^{1/3+\epsilon/9}$. Using $\tau(a) = o(1)$ as $a \to \infty$ and Lemma 2.3 we have $\sum \tau(a)/a \leq y^{-\epsilon/9 + o(1)}$, which is sufficient for our needs. So now assume that $a > y^{3-\epsilon}$. If also $\tau(a) \leq \exp(\epsilon \sqrt{\log a \log_2 a})$, a brief calculation with Lemma 2.6 suffices.

Now assume that $y^{3-\epsilon} < a \leq Z$ and $\tau(a) > \exp(\epsilon \sqrt{\log a \log_2 a})$. We forget at this point that $a$ is running over primitive nondeficient integers. For each fixed positive integer $j$,

$$(3) \quad \sum_{k \leq T} \frac{\tau(k)^j}{k} \leq \prod_{p \leq T} \left(1 + \frac{2^j}{p} + O_j(1/p^2)\right) \ll_{j} (\log T)^{2j}.$$
Let $T \geq y^{3-\epsilon}$ and consider values of $k \in (T, T^2]$ with $\tau(k) > z(T) := \exp(\epsilon \sqrt{\log T \log_2 T})$. We have
\[
\sum_{T < k \leq T^2 \atop \tau(k) > z(T)} \frac{\tau(k)}{k} \leq \frac{1}{z(T)^{j-1}} \sum_{T < k \leq T^2} \frac{\tau(k)^j}{k} \ll \frac{(\log T)^{2j}}{z(T)^{j-1}}.
\]
Thus, letting $T$ range over numbers of the form $(y^{3-\epsilon})^i$, for $i = 0, 1, 2, \ldots$,
\[
\sum_{a \geq y^{3-\epsilon} \atop \tau(a) > z(a)} \frac{\tau(a)}{a} \ll \frac{(\log y)^{2j}}{z(y^{3-\epsilon})^{j-1}}.
\]
If we choose $j = \lfloor \epsilon^{-1} \rfloor$, and let $\epsilon$ be arbitrarily small, this estimate completes the argument in the case $a \leq Z$.
Now we assume $a > Z$. Since $a \mid s(n)$ and $s(n) \leq n$, we have $a \leq x$. Consider the “$Z$-prefix” divisor $b$ of $a$. This is the largest divisor of $a$ in $[1, Z]$ with each prime dividing $b$ bounded above by each prime dividing $a/b$. We may assume that
\[
(4) \quad b > Z^{1/3}.
\]
Indeed, if $b \leq Z^{1/3}$, then every prime factor of $a/b$ is greater than $Z^{2/3}$, so that
\[
\frac{\sigma(a/b)}{a/b} = 1 + O\left(\frac{\log x}{Z^{2/3}}\right),
\]
since $a/b$ has fewer than $\log x$ distinct prime factors. Also, since $b$ is deficient, we have
\[
\frac{\sigma(b)}{b} \leq 2 - \frac{1}{b} \leq 2 - \frac{1}{Z^{1/3}}.
\]
The last two displays imply that for $x$ sufficiently large, we have $a$ deficient, a contradiction. Thus, we have (4).
Let $\mathcal{B}$ denote the set of integers $b \in (Z^{1/3}, Z]$ which are the $Z$-prefix for some primitive nondeficient $a \in (Z, x]$. We will show that there are very few integers $n \leq x$ with $s(n)$ divisible by some member of $\mathcal{B}$. To do this, we show that the set $\mathcal{B}$ is rather sparse.
Let $Y := Z^{1/\log_2 x}$, and assume that $b \in \mathcal{B}$ and $P(b) > Y$. Then, as above,
\[
\frac{\sigma(a/b)}{a/b} = 1 + O\left(\frac{\log x}{Y}\right),
\]
so for $a$ to be nondeficient, there is some absolute constant $c > 0$ such that
\[
2 > \frac{\sigma(b)}{b} \geq 2 - \frac{c \log x}{Y}.
\]
By Lemma 2.7, for $t \in (Z^{1/3}, Z]$, the number of integers $b \leq t$ which satisfy the last display is at most $t/\exp(\frac{1}{5} \sqrt{\log x / \log_3 x})$ for all sufficiently large $x$.
We also wish to count the number of $b \in \mathcal{B}$, $b \leq t$ with $P(b) \leq Y$. For this we ignore the condition $b \in \mathcal{B}$ and use only a standard result in [10]: the number of such $b$ is at most $t/(\log x)^{2 \log_3 x}$ for all sufficiently large $x$.
As in the case $a \leq Z$, we find that the number of integers $n \leq x$ with $b \mid s(n)$ for some $b \in \mathcal{B}$ (and assuming $n = pm$ with $p > x^{1/\log_3 x}$ and $p \nmid m$) is
\[
\ll x \log_3 x \sum_{b \in \mathcal{B}} \frac{\tau(b)}{b}.
\]
Again as before we treat the cases of $\tau(b)$ small and large separately. If $\tau(b) \leq (\log x)^{1/\log_3 x}$, then our estimates for the number of $b \leq t$ suffice to give an acceptable estimate. If
\[ \tau(b) > (\log x)^{\frac{1}{3} \log \log x}, \text{ we use (3) with } j = 2, \text{ getting} \]
\[ \sum_{Z^{1/3} < b \leq Z} \frac{\tau(b)}{b} \leq \frac{1}{(\log x)^{\frac{2}{3} \log \log x}} \sum_{Z^{1/3} < b \leq Z} \frac{\tau(b)^2}{b} \ll (\log x)^{4 - \frac{1}{8} \log \log x}. \]

This too is sufficiently small, so completing our proof. \( \square \)

2.2. Lower bounds. Here we prove Theorems 1.2, 1.3.

**Proof of Theorem 1.2.** Let \( x \) be large, let \( k \) be the least integer with \( 2^k > \log \log x \), and let \( p \) be a prime with \( 2^k < p < 2^{k+1} - 1 \). We shall consider numbers \( n = 2^kp^\alpha m \) satisfying (ii) and (iii). In fact, this is the correct order of magnitude for this count, since (iii) by itself already implies an upper bound of the same form. Since \( x^{1/6} < m \leq x^{1/3} \) and \( q \) is prime with \( q > x^{1/3} \). Since \( 2^kp \) is abundant, such numbers \( n \) are abundant. With no further restrictions on the parameters \( m, q \), the number of \( n \)‘s created in this way is \( \gg x / \log_2 x \log x \). We will show that with some mild restrictions on \( m, q \) we can be assured that \( s(n) \) is deficient, and these restrictions will only cut down on our count by a factor of \((\log_3 x)^2\).

We shall want the numbers \( m \) to satisfy:

(i) \( 2^k | \sigma(m) \),
(ii) \( p \nmid \sigma(m) \),
(iii) \( m \) has no prime factors below \((\log_3 x)^2\).

We can ensure that \( m \) satisfies (ii) and (iii) by eliminating those \( m \) divisible by any of:

- a proper prime power \( r^\alpha \geq p/2 \),
- a prime \( r \equiv -1 \pmod{p} \),
- a prime \( r \leq (\log_2 x)^2 \).

By Brun’s sieve [26, Theorem 2.2], there is an absolute constant \( \eta > 0 \) such that the number of \( m \leq t \) divisible by no prime \( r \leq t^{\eta} \) in the second or third bullet points is

\[ \gg t \prod_{r \leq (\log_2 x)^2} \left( 1 - \frac{1}{r} \right) \prod_{r \equiv -1 \pmod{p}} \left( 1 - \frac{1}{r} \right) \gg \frac{t}{\log_3 x} \exp \left( - \sum_{r \leq t^{\eta} \pmod{p}} \frac{1}{r} \right) \gg \frac{t}{\log_3 x}. \]

(The sum on \( r \) here is \( O(1) \), by the Brun–Titchmarsh inequality and the bound \( p \gg \log_2 t \).) The number of \( m \leq t \) divisible by some prime \( r > t^{\eta} \) with \( r \equiv -1 \pmod{p} \) is \( \ll t \sum_{r > t^{\eta} \pmod{p}} \frac{1}{r} \ll t/p \ll t/\log_2 x \), which is \( o(t/\log_3 x) \), as \( x \to \infty \). Finally, the number of \( m \leq t \) divisible by some proper prime power \( r^\alpha \geq p/2 \) is \( \ll t/\sqrt{p} \ll t/\sqrt{\log_2 x} \), and this is also \( o(t/\log_3 x) \). Thus, there are \( \gg t/\log_3 x \) values of \( m \leq t \) satisfying (ii) and (iii). In fact, this is the correct order of magnitude for this count, since (iii) by itself already implies an upper bound of the same form.

We now turn to (i). We apply the Turán–Kubilius inequality [42, Theorem 1, p. 302] to investigate the normal order of the additive function \( v_2(\sigma(m)) \). Uniformly for \( x^{1/6} < t \leq x^{1/3} \),

\[ \sum_{p^v \leq t} v_2(\sigma(p^v))p^{-v}(1 - p^{-1}) \sim 2 \log \log x \]

and

\[ \sum_{p^v \leq t} v_2(\sigma(p^v))p^{-v} \ll \log \log x. \]

Hence, the number of positive integers \( m \leq t \) where \( v_2(\sigma(m)) < \log \log x \) is \( O(t/\log_2 x) \). Since \( k \ll \log_3 x \), this is also an upper bound on the number of failures of (i).

We conclude that there are \( \asymp t/\log_3 x \) values of \( m \leq t \) satisfying all of (i)–(iii). By partial summation,

\[ (5) \sum_{m} \frac{1}{m} \sim \frac{\log x}{\log_3 x}, \]

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where the sum is over \( m \in (x^{1/6}, x^{1/3}] \) satisfying (i)–(iii).

For each \( m \), let \( f_0(m) = \sum_{r|m} \frac{1}{r} \), where \( r \) runs over primes. Using that \( m \) is free of prime factors up to \((\log_2 x)^2\) gives

\[
\sum_m \frac{1}{m} f_0(m) = \sum_{r > (\log_2 x)^2} \frac{1}{r} \sum_{m | r} \frac{1}{m} \leq \sum_{r > (\log_2 x)^2} \frac{1}{r^2} \prod_{p \leq x \atop p > (\log_2 x)^2} (1 - 1/p)^{-1},
\]
which is \( \ll \frac{\log x}{(\log_2 x)^2(\log_3 x)^2} \). In what follows, we will assume that

\[
(6) \quad f_0(m) \leq \frac{1}{2(\log_2 x)^2}. \]

Note that discarding contrary values of \( m \) changes \( \sum_m \frac{1}{m} \) only by \( O(\log x/(\log_3 x)^2) \) and so preserves the validity of (5).

We now complete the construction of a number \( n = 2^k pmq \) by choosing a prime \( q \in (x^{1/3}, x/2^k pm] \). Suppose an odd prime \( r \) divides \( s(n) \). Then

\[
s(n) = qs(2^k pm) + \sigma(2^k pm) \equiv 0 \pmod{r}.
\]

If \( r \mid s(2^k pm) \), then \( r \mid \sigma(2^k pm) \), so that \( r \mid 2^k pm \). Such a prime \( r \) cannot be 2 by assumption, and it cannot be \( p \), since \( p \mid \sigma(2^k) \) and \( p \mid \sigma(m) \). Thus, \( r \mid m \), and so \( r > (\log_2 x)^2 \). We conclude that if \( 2 < r \leq (\log_2 x)^2 \) and \( r \mid s(n) \), then \( r \mid s(2^k pm) \).

We now restrict the \( q's \) that are available to us by removing for each odd prime \( r \leq (\log_2 x)^2 \) a single coprime residue class that will ensure that \( r \mid s(n) \). By another application of Brun’s sieve, this cuts down the number of choices for \( q \) by a factor proportional to \( \log_3 x \).

For a surviving prime \( q \), let

\[
f_1(q) = \sum_{2 < r < x^{1/4} \atop r \mid m \atop r \mid s(2^k pm)} \frac{1}{r}.
\]

If \( r \mid s(2^k pmq) \) and \( r \nmid m \), then as we have seen, \( q \) is in a fixed residue class modulo \( r \). Using that \( q \) also avoids one residue class modulo each prime not exceeding \((\log_2 x)^2\), the Brun–Titchmarsh inequality implies that the average value of \( f_1(q) \), over surviving primes \( q \), is \( \ll 1/(\log_3 x)^2 \log_2 x \). Keeping (6) in mind, we conclude that there are \( \gg x \log_2 x / 2^{k} pm m x \log_2 x \log_3 x \) values of \( q \) for which

\[
(7) \quad \sum_{r > 2 \atop r \mid s(2^k pmq)} \frac{1}{r} \leq \frac{1}{(\log_2 x)^2}.
\]

We finally note that for \( m, q \) chosen as above, we have \( 2^{k+1} \mid \sigma(2^k pmq) \), so that \( 2^k \mid s(2^k pmq) \). Hence, \( s(2^k pmq)/2^k \) has only odd prime factors. From (7), we conclude that \( s(2^k pmq) \) is deficient. The theorem then follows with the estimate:

\[
\sum_p \sum_m \sum_q 1 \gg \sum_p \sum_m 2^k pm \frac{x}{\log_2 x \log_3 x} \sum_p \frac{x}{2^{k} pm (\log_3 x)^2} \gg \frac{x}{2^{k} (\log_3 x)^3} \gg \frac{x}{\log_2 x (\log_3 x)^3},
\]

where we have used (5).

\( \square \)

**Proof of Theorem 1.3.** We take \( k \) as the least integer with \( 2^k > 10 \log \log x \). We consider numbers \( n = 2^k mqr \leq x \) with \( q, r \) primes, \( x^{1/8} < m \leq x^{1/4}, x^{1/4} < q \leq x^{1/3} \), and \( r > x^{1/3} \).

As in the last proof, we impose some conditions on \( m \). Namely, we insist that \( m \) satisfy all of

- (i’’) \( m \) has no prime factors below \( (\log_2 x)^2 \),
- (ii’’) \( 2^k \mid \sigma(m) \),
(iii') (6) holds, 
(iv') for $I := [2^k, 2^{k+1} - 1]$, we have $\sum_{p \in I, p \mid \sigma(m)} \frac{1}{p} > \frac{1}{10} \frac{1}{\log_3 x}$.

We will show we can construct a collection of such $m$ that satisfies

$$\sum_m \frac{1}{m} \geq \frac{\log x}{\log_3 x}.$$  

(8)

It will turn out that the most stringent restriction is the first one.

By an elementary sieve, those $m \in (x^{1/8}, x^{1/4}]$ satisfying (i') have

$$\sum_m \frac{1}{m} \sim \frac{1}{16e^7 \log_3 x},$$

as $x \to \infty$. By arguments seen above, restricting $m$ further to satisfy (ii') changes the value of $\sum_m \frac{1}{m}$ by $O(\log x/\log_2 x)$, and restricting $m$ to satisfy (iii') affects the sum only by $O(\log x/(\log_3 x)^2)$. Hence, (9) still holds if $m$ is restricted to satisfy all of (i')–(iii').

We now turn to (iv'). Changing $\sum_m \frac{1}{m}$ by $O(\log x/\sqrt{\log_2 x})$, we can assume that $m$ is not divisible by any proper prime power exceeding $5 \log_2 x$. If $p \in I$ and $p \mid \sigma(m)$, then there is a prime power $\ell^\alpha \mid m$ with $p \mid \sigma(\ell^\alpha)$; since $\ell^\alpha > p/2 > 5 \log_2 x$, it follows that $\alpha = 1$. Let

$$f_2(m) = \sum_{p \in I \atop p \mid \sigma(m)} \frac{1}{p}.$$  

Using that $m \leq x^{1/4}$ and that $m$ is free of prime factors up to $(\log_2 x)^2$, we see that as $x \to \infty$,

$$\sum_m \frac{1}{m} f_2(m) \leq \sum_{p \in I} \frac{1}{p} \sum_{\ell \equiv 1 \pmod{m}} \sum_{m: \ell \mid m} \frac{1}{m} \leq \left( \frac{1}{8} + o(1) \right) \frac{\log x}{\log_3 x} \sum_{p \in I} \frac{1}{p} \sum_{\ell \leq x} \frac{1}{\ell}.$$  

From [35, Theorem 1 and Remark 1], the final sum on $\ell$ is at most $(1 + o(1)) \frac{\log \log x}{p}$, as $x \to \infty$, uniformly for $p \in I$. Inserting this estimate above,

$$\sum_m \frac{1}{m} f_2(m) \leq \left( \frac{1}{8} + o(1) \right) \frac{\log x}{\log_3 x} \log_2 x \cdot \sum_{p \in I} \frac{1}{p^2} \leq \left( \frac{1}{80} + o(1) \right) \frac{\log x}{(\log_3 x)^2}.$$  

Hence, the contribution to the left-hand side of (9) from those values of $m$ with $f_2(m) \geq \frac{1}{2 \log_3 x}$ is at most $(\frac{1}{40} + o(1)) \frac{\log x}{\log_3 x}$. Since $\frac{\log x}{\log_3 x} - \frac{1}{40} > \frac{1}{100}$, after discarding these $m$ we may still assert (8). Moreover, since $\sum_{p \in I} \frac{1}{p} \sim \frac{\log 2}{\log_2 x}$ and $\log 2 - \frac{1}{2} > \frac{1}{10}$, we have

$$\sum_{p \in I \atop p \mid \sigma(m)} \frac{1}{p} > \frac{1}{10 \log_3 x}$$

for all remaining $m$ (and all large $x$).

Note that for any choice of primes $q, r$ greater than $m$, the number $2^k m qr$ is deficient. We shall show that with some mild restrictions on $q, r$, we will have $s(2^k m qr)$ abundant. For $m$ satisfying the above restrictions, let $S(m)$ be the set of primes $p \in I$ with $p \mid \sigma(m)$ and let $S_m$ be the product of the members of $S(m)$.

We next consider primes $q \in (x^{1/4}, x^{1/3}]$. They are chosen with the requirements that both $q + 1$ and $s(2^k m q)$ are coprime to $S_m$. We have $s(2^k m q) = qs(2^k m) + \sigma(2^k m)$. By our choices of $k, m$, we have $\sigma(2^k m)$ coprime to $S_m$. So for a prime $p \mid \gcd(s(2^k m q), S_m)$, $q$ is in a particular residue class $p$ mod $p$. Similarly this holds too for $p \mid q + 1$. Thus, the density of the primes $q$ which avoid these classes is at least

$$\prod_{p \mid S_m} \frac{p - 3}{p - 1} \sim 1, \quad \text{as } x \to \infty.$$
So, almost all of the primes \( q \in (x^{1/4}, x^{1/3}] \) may be used, and for these primes \( q \), we have \( \gcd(s(2^k mq), S_m) = 1 \) and \( \gcd(q + 1, S_m) = 1 \).

We now choose primes \( r \in (x^{1/3}, x/(2^k mq)] \) with \( \gcd(s(2^k mqr), S_m) > 1 \). For such a prime \( r \), \( s(2^k mqr) \) is abundant. Let \( p \mid S_m \). Note that for any prime \( r \mid 2^k mq \), we have \( s(2^k mqr) = rs(2^k mq) + \sigma(2^k mq) \). Further, we have \( p \nmid s(2^k mq) \) and \( p \nmid \sigma(2^k mq) \).

By the prime number theorem for progressions, the fraction of \( r \) with \( p \mid s(2^k mqr) \) is asymptotically \( 1/(p - 1) \). Similarly, if \( pp' \mid S_m \), the fraction of choices for \( r \) with \( pp' \mid s(2^k mqr) \) is asymptotically \( O(1/pp') \). Since

\[
\sum_{p \mid S_m} \frac{1}{p - 1} \gg \frac{1}{\log x}, \quad \sum_{pp' \mid S_m} \frac{1}{pp'} \ll \frac{1}{(\log x)^2},
\]

we have that the proportion of primes \( r \) with \( \gcd(s(2^k mqr), S_m) > 1 \) is \( \gg 1/\log x \).

Putting these thoughts together, we have that the number of choices for \( m, q, r \)

\[
\sum_m \sum_q \sum_r 1 \gg \sum_m \sum_q 2^{k m q} x \log x \log_3 x \gg \sum_m x 2^{k m \log x \log_3 x} \gg \frac{x}{2^k (\log x)^2}.
\]

Since \( 2^k \ll \log \log x \), this completes the proof of the theorem. \( \square \)

3. The density of nonaliquot numbers

3.1. A heuristic argument. For a positive integer \( y \), let

\[ A_y = \text{lcm}[1, 2, \ldots, y]. \]

For a positive integer \( a \mid A_y \), let

\[ T_a = \{ n : \gcd(n, A_y) = a \}, \]

and let \( T_a(x) \) denote the set of members of \( T_a \cap [1, x] \). We assume that \( s \) maps \( T_a \) to \( T_a \) for each \( a \mid A_y \). This is asymptotically true if \( n > e^{\alpha y} \) and \( y \to \infty \). Further, we assume that for \( n \in T_a \) we have \( \sigma(n)/n \approx \sigma(a)/a \). This is asymptotically true up to sets of vanishing density as \( y \to \infty \). We assume in addition that \( s \) is a random map subject to the prior assumptions.

If \( s(m) \in T_a(x) \), then \( m \in T_a \) and \( m \leq x/(\sigma(a)/a - 1) = (a/s(a))x \), so that \( m \in T_a((a/s(a))x) \). The probability that an integer \( n \in T_a(x) \) is not in the range of \( s \) is about

\[ \left(1 - \frac{1}{|T_a((a/s(a))x)|}\right)^{|T_a((a/s(a))x)|}. \]

Now \( |T_a(x)| \sim (\varphi(A_y)/a A_y)x \), so the above probability is asymptotically

\[ \exp\left(-\frac{a}{s(a)}\right). \]

We thus would have that the density \( \Delta \) of even numbers not in the range of \( s \) is

\[
\Delta = \lim_{y \to \infty} \frac{\varphi(A_y)}{A_y} \sum_{a \mid A_y}^* \frac{1}{a} e^{-a/s(a)},
\]

where the * indicates that \( a \) runs over even numbers. Since almost all odd numbers are in the image of \( s \), \( \Delta \) is also the density of nonaliquot numbers.

We now prove that the limit (10) exists. Let

\[
P_y = \frac{\varphi(A_y)}{A_y} = \prod_{p \leq y} \left(1 - \frac{1}{p}\right), \quad S_y = \sum_{P(a) \leq y}^* \frac{1}{a} e^{-a/s(a)}, \quad \bar{S}_y = \sum_{a \mid A_y}^* \frac{1}{a} e^{-a/s(a)}.
\]
Note that $S'_y \leq S_y$. The difference comes from those $a$ with $P(a) \leq y$ and $a \nmid A_y$. Such a number $a$ is divisible by $p^k$ for some prime $p^k > y$ and $p \leq y$. Writing $a = p^ka'$, we have

$$\sum_{a|A_y}^* \frac{1}{a} e^{-a/s(a)} \leq \sum_{p \leq y} \sum_{k \geq \log y \log p}^* \frac{1}{p^k} \sum_{a|A_y}^* \frac{1}{a} \ll \log y \sum_{p \leq y}^* \frac{1}{y} \ll 1,$$

Since, by Mertens’ theorem, $P_y \sim e^{-\gamma}/\log y$ as $y \to \infty$, we have

$$P_y S_y - P_y S'_y \ll \frac{1}{\log y}.$$  

Thus the limit in (10) exists if and only if $\lim_{y \to \infty} P_y S_y$ exists, in which case this latter limit is also $\Delta$.

For a prime $q$, we have $P_q = P_{q-1}(1 - \frac{1}{q})$ and, for $q > 2$,

$$S_q - S_{q-1} = \sum_{a_q \leq q}^* \frac{1}{a_q} e^{-a_q/s(a_q)} > \sum_{a_q \leq q}^* \frac{1}{a_q} e^{-a_q/s(a_q)} = \frac{1}{q-1} S_{q-1}.$$  

Thus,

$$P_q S_q > P_{q-1} \left(1 - \frac{1}{q}\right) S_{q-1} \left(1 + \frac{1}{q-1}\right) = P_{q-1} S_{q-1}.$$  

Since $P_y S_y$ is always at most $\frac{1}{2}$, it follows that the limit in (10) exists and

$$\Delta = \lim_{y \to \infty} e^{-\gamma} \log y \sum_{a \leq y, a|A_y}^* \frac{1}{a} e^{-a/s(a)} = \lim_{y \to \infty} \frac{e^{-\gamma}}{\log y} \sum_{P(a) \leq y}^* \frac{1}{a} e^{-a/s(a)}.$$  

We would like to compute a numerical approximation to $\Delta$, and it may seem that the first expression in (12) is more easily computed than the second, since it involves a finite sum while the latter involves an infinite sum. However, we claim the infinite sum can be replaced with the sum over $a \leq y$, while changing $e^{-\gamma}$ to 1. That is,

$$\Delta = \lim_{y \to \infty} \frac{1}{\log y} \sum_{a \leq y}^* \frac{1}{a} e^{-a/s(a)}.$$  

The equation (13) may be proved as follows: Introduce a parameter $z \geq 2$, which will eventually be chosen as a slow-growing function of $y$. Define $h_z(n) = \sum_{d|n, 1 < d \leq z} 1/d$. For even $a$,

$$0 \leq e^{-\frac{a}{\pi(a)}} - e^{-\frac{1}{h_z(a)}} = e^{-\frac{a}{\pi(a)}} \left(1 - e^{\frac{1}{\pi(a)/a} - \frac{1}{h_z(a)}}\right) \ll e^{-\frac{a}{\pi(a)}} \left(\frac{s(a)}{a} - h_z(a)\right).$$

Thus,

$$0 \leq \sum_{a \leq y}^* \frac{1}{a} e^{-\frac{a}{\pi(a)}} - \sum_{a \leq y}^* \frac{1}{a} e^{-\frac{1}{h_z(a)}} \ll \sum_{a \leq y}^* \frac{1}{a} e^{-\frac{a}{\pi(a)}} \sum_{d|a} \frac{1}{d} \sum_{d > z} \frac{1}{d}$$

$$\leq \sum_{d > z} \frac{1}{d} \sum_{a \leq y}^* \frac{1}{a} \leq \log y \sum_{d > z} \frac{1}{d^2} \ll \frac{1}{z} \log y.$$  

A nearly identical computation gives

$$\sum_{P(a) \leq y}^* \frac{1}{a} e^{-a/s(a)} - \sum_{P(a) \leq y}^* \frac{1}{a} e^{-\frac{1}{h_z(a)}} \ll \frac{1}{z} \log y.$$  

$$\sum_{P(a) \leq y}^* \frac{1}{a} e^{-a/s(a)} - \sum_{P(a) \leq y}^* \frac{1}{a} e^{-\frac{1}{h_z(a)}} \ll \frac{1}{z} \log y.$$
Note that \( h_z(a) = h_z(\gcd(a, A_z)) \), and that when \( a \) is even, \( d = \gcd(a, A_z) \) is also even. Hence, for all large \( y \) we have

\[
\sum_{a \leq y} \frac{1}{a} e^{-\frac{1}{\pi_z(a)}} = \sum_{d | A_z} \sum_{P(a) \leq y \atop \gcd(a, A_z) = d} \frac{1}{a} e^{-\frac{1}{\pi_z(a)}} = \sum_{d | A_z} \frac{1}{a} \sum_{\mu(j)} = \sum_{d | A_z} \frac{1}{a} \sum_{j | a/d} \frac{1}{j} \sum_{P(a) \leq y} \frac{1}{j}.
\]

Again almost identically, we have

\[
\sum_{a \leq y} \frac{1}{a} e^{-\frac{1}{\pi_z(a)}} = \sum_{d | A_z} \frac{1}{d} e^{-\frac{1}{\pi_z(d)}} \sum_{a \leq y/d} \frac{1}{a}.
\]

By Mertens’ theorem (with the elementary error estimate shown in [42, Theorem 10, p. 17]),

\[
\frac{1}{\log y} \sum_{P(a) \leq y} \frac{1}{a} = \frac{e^\gamma}{\log y} \sum_{a \leq y} \frac{1}{a} + O \left( \frac{1}{\log y} \right) = \frac{e^\gamma}{\log y} \sum_{a \leq y/d} \frac{1}{a} + O \left( \frac{\log(dj) + 1}{\log y} \right).
\]

Making this substitution, we find that

\[
\left| e^{-\gamma} \frac{\log y}{\sum_{P(a) \leq y} \frac{1}{a} e^{-\frac{1}{\pi_z(a)}}} - \frac{1}{\log y} \sum_{a \leq y} \frac{1}{a} e^{-\frac{1}{\pi_z(a)}} \right| \ll \frac{1}{\log y} \sum_{d | A_z} \frac{\log dj}{dj} \ll \frac{(\log z)^3}{\log y}.
\]

By choosing \( z = \exp((\log y)^{1/4}) \), say, from (11), (14), (15), and (16), we see that

\[
\left| e^{-\gamma} \frac{\log y}{\sum_{a | A_y} \frac{1}{a} e^{-a/s(a)}} - \frac{1}{\log y} \sum_{a \leq y} \frac{1}{a} e^{-a/s(a)} \right| \ll \frac{1}{(\log y)^{1/4}}.
\]

Thus, (13) follows. This completes the heuristic justification of Conjecture 1.4.

### 3.2. Computations

The following strategy was used to construct Table 1. Explicit computations on Goldbach’s conjecture have shown that every even number in \([8,4 \cdot 10^{18}]\) is a sum of two distinct primes \( p \) and \( q \) (see [30]). Hence, 5 is the only odd nonaliquot number up to \( 10^{10} \). In [38, §5], an algorithm is described that enumerates the elements of \( (N)^/[1, N] \) in time \( \tilde{O}(N) \). We ran the algorithm to \( N = 10^{10} \). We made one modification, necessitated by memory constraints: Instead of using a 0-1 array to record which numbers had appeared so far, the even numbers in the range of \( s \) were written to disk as they were found by the algorithm, and postsorted later with the GNU sort utility. The final data file required 36.1 GB.

On the computation of an approximation for \( \Delta \), this is straightforward using (13) for large values of \( y \). Note that \( \Delta \) is effectively computable, since the above calculations show that for each fixed integral value of \( z \geq 2 \),

\[
\lim_{y \to \infty} e^{-\gamma} \frac{\log y}{\sum_{d | A_z} \frac{\mu(j)}{dj} e^{-\frac{1}{\pi_z(d)}} \sum_{P(a) \leq y} \frac{1}{a} e^{-\frac{1}{\pi_z(a)}}} = \sum_{d | A_z} \frac{\mu(j)}{dj} e^{-\frac{1}{\pi_z(d)}}
\]

is within \( 4/z \) of \( \Delta \). We have not used this approach, but we record some intermediate values for the expression in (13) at \( y = 10^j \) and \( 2 \cdot 10^j \) for \( 6 \leq j \leq 10 \):

- \( 10^6 : 0.166923 \), \( 2 \cdot 10^6 : 0.167484 \), \( 10^7 : 0.168599 \), \( 2 \cdot 10^7 : 0.169014 \), \( 10^8 : 0.169857 \), \( 2 \cdot 10^8 : 0.170176 \), \( 10^9 : 0.170834 \), \( 2 \cdot 10^9 : 0.171088 \), \( 10^{10} : 0.171617 \), \( 2 \cdot 10^{10} : 0.171822 \).
Lemma says that we have a strong bound on the number of solutions to $\sigma(n)$ one would decide if a given very large number is of the required form. The evidence is somewhat less compelling than with the nonaliquots. At least the count densities seem to have peaked and are now slowly descending, perhaps to the limit $\Delta_{\varphi}$. With both the nonaliquots and the $\varphi$ analogues, the heuristic is very much an “at infinity” type of argument, while the counts are definitely very finite. It might be of some interest to try to do a statistical sample of random numbers that are much larger than our exhaustive counts, but it is not clear how one would decide if a given very large number $n$ is of the required form.

Here are the intermediate calculations of the expression in (17) for various values of $y$:

- $10^6$: 0.089193
- $2 \cdot 10^6$: 0.089353
- $10^7$: 0.089673
- $2 \cdot 10^7$: 0.089791
- $10^8$: 0.090033
- $2 \cdot 10^8$: 0.090124
- $10^9$: 0.090312
- $2 \cdot 10^9$: 0.090385
- $10^{10}$: 0.090536
- $2 \cdot 10^{10}$: 0.090595

### 4. Primitive friendly pairs

We begin by collecting some results needed for the proof of Theorem 1.5. The following lemma says that we have a strong bound on the number of solutions to $\sigma(n)/n = \alpha$ even if we specify only the denominator of $\alpha$. For the proof of this result, in a more precise form, see [31, Theorem 4.1].

**Lemma 4.1.** The number of $n \leq x$ for which $\sigma(n)/n$ can be written as a fraction with a given denominator $d$ is at most $x^{\alpha(1)}$, as $x \to \infty$, uniformly in $d \leq x$. 

<table>
<thead>
<tr>
<th>$x$</th>
<th>$U_{\varphi}(x)$</th>
<th>$D_{\varphi}(x)$</th>
<th>$x$</th>
<th>$U_{\varphi}(x)$</th>
<th>$D_{\varphi}(x)$</th>
</tr>
</thead>
<tbody>
<tr>
<td>10000000</td>
<td>11355049</td>
<td>0.1136</td>
<td>1000000000</td>
<td>113482572</td>
<td>0.1135</td>
</tr>
<tr>
<td>20000000</td>
<td>22718595</td>
<td>0.1136</td>
<td>2000000000</td>
<td>226723208</td>
<td>0.1134</td>
</tr>
<tr>
<td>30000000</td>
<td>34079552</td>
<td>0.1136</td>
<td>3000000000</td>
<td>339811983</td>
<td>0.1133</td>
</tr>
<tr>
<td>40000000</td>
<td>45433178</td>
<td>0.1136</td>
<td>4000000000</td>
<td>452815609</td>
<td>0.1132</td>
</tr>
<tr>
<td>50000000</td>
<td>56783449</td>
<td>0.1136</td>
<td>5000000000</td>
<td>565739195</td>
<td>0.1131</td>
</tr>
<tr>
<td>60000000</td>
<td>68129835</td>
<td>0.1136</td>
<td>6000000000</td>
<td>678597689</td>
<td>0.1131</td>
</tr>
<tr>
<td>70000000</td>
<td>79472816</td>
<td>0.1136</td>
<td>7000000000</td>
<td>791405364</td>
<td>0.1131</td>
</tr>
<tr>
<td>80000000</td>
<td>90811869</td>
<td>0.1136</td>
<td>8000000000</td>
<td>904168542</td>
<td>0.1130</td>
</tr>
<tr>
<td>90000000</td>
<td>102146368</td>
<td>0.1136</td>
<td>9000000000</td>
<td>1016908189</td>
<td>0.1130</td>
</tr>
<tr>
<td>1000000000</td>
<td>1129598504</td>
<td>0.1130</td>
<td>10000000000</td>
<td>1129598504</td>
<td>0.1130</td>
</tr>
</tbody>
</table>

Table 4. Counts $U_{\varphi}(x)$ of numbers not in the range of $s_{\varphi}(\cdot)$ to various heights $x$ from $10^8$ through $10^{10}$, along with the corresponding densities $D_{\varphi}(x)$. 

### 3.3. A $\varphi$ analogue.

Let $s_{\varphi}(n) = n - \varphi(n)$. As with $s(n)$, partial results on Goldbach’s conjecture imply that the range of $s_{\varphi}(\cdot)$ contains all odd numbers but for a set of asymptotic density 0. Little is known about even values. In [19] Erdös asks if a positive proportion of even numbers are of the form $s_{\varphi}(n)$ and if a positive proportion of even numbers are not of the form $s_{\varphi}(n)$. The proof in [29] does give a positive proportion of evens that are of the form $s_{\varphi}(n)$, but the even numbers not of this form remain an enigma. However, as with $s(n)$ the problem is amenable to calculation and heuristics. Let $U_{\varphi}(x)$ denote the number of integers in $[1, x]$ not of the form $s_{\varphi}(n)$. By an analogous algorithm from [38], we have extended the calculations of that paper 100-fold, see Table 4. The data file had 42.1 GB. (This file is larger than the one with nonaliquots since there are more even numbers $n$ with $s_{\varphi}(n) \leq 10^{10}$ than there are with $s(n) \leq 10^{10}$.)

The conjectural density is

$$\Delta_{\varphi} := \lim_{y \to \infty} \frac{1}{\log y} \sum_{a \leq y} \frac{1}{a} e^{-a/s_{\varphi}(a)}.$$ 

At $y = 2 \cdot 10^{10}$ this expression rounds to 0.090595. It is difficult to say if the actual counts are supportive of the conjectural density $\Delta_{\varphi}$; the evidence is somewhat less compelling than with the nonaliquots. At least the count densities seem to have peaked and are now slowly descending, perhaps to the limit $\Delta_{\varphi}$. With both the nonaliquots and the $\varphi$ analogues, the heuristic is very much an “at infinity” type of argument, while the counts are definitely very finite. It might be of some interest to try to do a statistical sample of random numbers that are much larger than our exhaustive counts, but it is not clear how one would decide if a given very large number $n$ is of the required form.
The next lemma is implicit in the proof of [23, Theorem 11]; for an explicit statement and proof, see [31, Lemma 4.2].

**Lemma 4.2.** Let \( m \) be an integer in \([1, x]\). The number of integers \( n \in [1, x] \) for which \( \text{rad}(n) \mid m \) is \( x^{o(1)} \), as \( x \to \infty \), uniformly in \( m \).

The next result, which was inspired by Erdős’s method of proof in [16], appears as [34, Lemma 3.1].

**Lemma 4.3.** Given a natural number \( m \), the following algorithm outputs a unitary divisor \( a \) of \( m \) with \( \gcd(a, \sigma(a)) = 1 \). Moreover, at most \( x^{o(1)} \) inputs \( m \leq x \) correspond to the same output \( a \), as \( x \to \infty \).

**Algorithm A:**

**Input:** A natural number \( m \)

**Output:** A divisor \( a \) of \( m \) for which \( a \mid m \) and \( \gcd(a, \sigma(a)) = 1 \)

Factor \( m = p_1^{e_1} p_2^{e_2} \cdots p_k^{e_k} \), where \( p_1 > p_2 > \cdots > p_k \).

\( a \leftarrow 1 \) \hspace{1cm} // Initialize

**for** \( i = 1 \) **to** \( k \) **do**

\( \text{if } \gcd(\sigma(p_i^{e_i} a), p_i^{e_i} a) = 1 \) **then**

\( a \leftarrow p_i^{e_i} a \)

end

end

**return** \( a \)

**Proof of Theorem 1.5.** Suppose that \( n \) and \( n' \) form a primitive friendly pair contained in \([1, x]\). Write \( n = N_1 N_2 \), where \( N_1 \) is squarefree, \( N_2 \) is squarefull, and \( \gcd(N_1, N_2) = 1 \). Define \( n_2 \) and \( n_3 \) by the equation \( N_2 = n_2 n_3 \), where \( n_2 \) is the output of Algorithm A when \( m = N_2 \). Note that \( n_2 \) and \( n_3 \) are coprime squarefull integers, since \( n_2 \) is a unitary divisor of \( N_2 \). Write

\[
\frac{\sigma(n)}{n} = \frac{\sigma(n')}{n'} = \frac{A}{B},
\]

and let \( U = \text{rad}(\gcd(B, n_2)) \). Define \( N_1', N_2', n_2', n_3', U' \) analogously, but now with reference to \( n' \) instead of \( n \). Since \( \gcd(n_2, n_2') = 1 \), it follows that \( \gcd(U, U') = 1 \). With \( L = \lfloor \log x \rfloor \), choose \( \theta_1, \theta_2, \theta_3, \delta \in \{ \frac{1}{L}, \frac{2}{L}, \ldots, \frac{L-1}{L} \} \) for which

\[
N_1' \in [e^{-x^{\theta_1}}, x^{\theta_1}], \quad n_2' \in [e^{-x^{\theta_2}}, x^{\theta_2}], \quad n_3' \in [e^{-x^{\theta_3}}, x^{\theta_3}], \quad U' \in [e^{-x^{\delta}}, x^{\delta}].
\]

Define \( \theta_1', \theta_2', \theta_3', \delta' \) analogously. Finally, choose \( \theta \in \{ \frac{1}{L}, \frac{2}{L}, \ldots, \frac{L-1}{L} \} \) with

\[
B \in [e^{-x^{\theta}}, x^{\theta}].
\]

The number of possibilities for the \( \theta_i \), the \( \theta'_i \), \( \delta \), \( \delta' \), and \( \delta_0 \) is at most \( L^9 = x^{o(1)} \). Thus, it is enough to prove that the number of pairs \( n, n' \) corresponding to a given choice of these variables is bounded by \( x^{1/2+o(1)} \), as \( x \to \infty \). We will establish four upper bounds on the number of these pairs \( n, n' \). The theorem will follow by combining these upper bounds.

Our first upper bound uses that

\[
B^2 \mid n_2 n_3 n_2' n_3', \tag{18}
\]

To prove this, write \( B = B_1 B_2 \), where \( B_1 \) is squarefree, \( B_2 \) is squarefull, and \( \gcd(B_1, B_2) = 1 \). Since \( B \) divides \( n \) and \( n' \), \( B_1 \mid n_2 n_3 \) and \( B_2 \mid n_2 n_3' \), so that \( B_2^2 \mid n_2 n_3 n_2' n_3' \). Now take a prime \( p \mid B_1 \). Since \( n \) and \( n' \) have no common unitary divisor, \( p \) cannot exactly divide both \( n \) and \( n' \). So either \( p^2 \mid n_2 n_3 \) or \( p^2 \mid n_2' n_3' \). Hence, \( B_1^2 \mid n_2 n_3 n_2' n_3' \) and (18) holds.

By Lemma 4.3, given \( n_2 \) there are only \( x^{o(1)} \) possibilities for \( n_3 \). Similarly, \( n_2' \) determines \( n_3' \) in at most \( x^{o(1)} \) ways. Thus, \( n_2 \) and \( n_2' \) determine the product \( n_2 n_3 n_2' n_3' \) up to \( x^{o(1)} \) possibilities. Now (18), along with the maximal order of the divisor function, shows that \( n_2 \) and \( n_2' \) determine
the pair \( n, n' \) in at most \( x^{o(1)} \) ways. Since the number of squarefull numbers in \([1, t]\) is \( O(t^{1/2}) \), the number of pairs is at most

\[
x^{\theta_2/2 + \theta'_2/2 + o(1)},
\]

as \( x \to \infty \).

For a second upper bound, expanding \( \frac{A}{B} = \frac{\sigma(N)}{N} \), we find that \( B\sigma(N_1)\sigma(n_2)\sigma(n_3) = AN_1n_2n_3 \). Since \( \gcd(n_2, \sigma(n_2)) = 1 \),

\[
n_2 \mid B\sigma(N_1)\sigma(n_3),
\]

so that

\[
\gcd(n_2) \mid U \cdot \sigma(N_1)\sigma(n_3).
\]

By Lemma 4.2, the number of possibilities for \( n = N_1n_2n_3 \) is at most

\[
x^{\delta + \theta_1 + \theta_2/2 + o(1)}.
\]

By Wirsing’s theorem (or Lemma 4.1), this is also a bound on the number of possibilities for the pair \( n, n' \). A symmetric argument gives our third upper bound,

\[
x^{\theta' + \theta_1'/2 + o(1)}.
\]

To prove the fourth upper bound, note that since \( \gcd(U, U') = 1 \), we have \( U^2U'^2 \mid B \).

Consequently, the number of possibilities for \( B \) is

\[
\ll \sum_{u_1 \geq e^{-1}x^{\delta}} \frac{x^\theta}{u_1^2u_2^{\theta'}} \ll x^{\theta - \delta - \delta'}.
\]

By Lemma 4.1, the number of possibilities for the pair \( n, n' \) is at most

\[
x^{\theta - \delta - \delta' + o(1)}.
\]

Taking the geometric mean of the four upper bounds (19), (20), (21), and (22), the number of possibilities for the pair \( n, n' \) does not exceed

\[
x^{\frac{1}{2}(\theta_1 + \theta'_1 + \frac{1}{2}(\theta_2 + \theta'_2 + \theta_3 + \theta'_3) + \theta) + o(1)}.
\]

Since \( n = N_1n_2n_3 \leq x \), we have \( \theta_1 \leq 1 - \theta_2 - \theta_3 + o(1) \); similarly, \( \theta'_1 \leq 1 - \theta'_2 - \theta'_3 + o(1) \).

From (18), \( 2\theta \leq \theta_2 + \theta_3 + \theta'_2 + \theta'_3 + o(1) \). Substituting these bounds for \( \theta, \theta_1 \) and \( \theta' \) into (23) gives a final upper bound of \( x^{\frac{1}{2} + o(1)} \), as \( x \to \infty \).

### 4.1. Computations.

To compile Table 2, for each \( n \leq 10^L \), we found all solutions \( m \leq 10^L \) with \( \frac{\sigma(m)}{m} = \frac{\sigma(n)}{n} \), recording only those solutions \( m \) with \( \gcd(m, n) = 1 \). To find the values of \( m \), we used the gp script solveBA of Michel Marcus (see [41]), which was based on an earlier program by Jan Munch Pedersen. The script solveBA uses a recursive algorithm to find all solutions \( n \) up to a given limit satisfying an equation of the form \( \frac{\sigma(n)}{n} = \frac{A}{B} \). The values of \( \gcd(m, n) \) were computed using Algorithm B3 below.

---

**Algorithm B1:**

**Input:** A natural number \( n \) and a divisor \( d \) of \( n \)

**Output:** The smallest \( D \mid n \) with \( d \mid D \)

\( D \leftarrow d \).

Iterate \( D \mapsto D \cdot (D, n/D) \) until \( D \) stabilizes.

**return** \( D \)

---

**Algorithm B2:**

**Input:** A natural number \( n \) and a divisor \( d \) of \( n \)

**Output:** The largest \( D \mid n \) with \( D \mid d \)

\( D \leftarrow d \).

Iterate \( D \mapsto D/(D, n/D) \) until \( D \) stabilizes.

**return** \( D \)
\section*{Algorithm B3:}

\begin{itemize}
\item \textbf{Input:} Natural numbers \( m \) and \( n \)
\item \textbf{Output:} \( \gcd(m, n) \)
\item \( d \leftarrow \gcd(m, n) \)
\item \( d_1 \leftarrow \) output of \( B_1 \) with input \( d, m \)
\item \( d_2 \leftarrow \) output of \( B_1 \) with input \( d, n \)
\item \( d_3 \leftarrow d_1 d_2 / d \)
\item \textbf{return} the output of \( B_2 \) with the pair \( d, d_3 \)
\end{itemize}

5. The distribution function for \( F(n; x) \)

Define a \textit{friendly} \( k \)-set as a set of \( k \) distinct integers all of which share the same value of \( \sigma(n)/n \), and call such a set \textit{primitive} if there is no \( d > 1 \) that is a unitary divisor of each element. The proof of Theorem 1.6 will use the following crude upper bound on the count of primitive friendly \( k \)-sets.

\textbf{Lemma 5.1.} Fix \( k \geq 2 \). The number of primitive friendly \( k \)-sets contained in \([1, x]\) is at most \( x^{1 - \frac{1}{k+1} + o(1)} \), as \( x \to \infty \).

When \( k = 2 \), Lemma 5.1 gives an upper bound of \( x^{2/3+o(1)} \), which is inferior to the result of Theorem 1.5. However, the proof of Lemma 5.1 is substantially simpler.

\textbf{Proof.} List the elements of the set as \( n_1, \ldots, n_k \). Write the common value of \( \sigma(n_i)/n_i \) as \( A/B \) in lowest terms. If \( \text{rad}(B) \leq x^{1 - \frac{1}{k+1}} \), then there are at most \( x^{1 - \frac{1}{k+1} + o(1)} \) possible values of \( B \), by Lemma 4.2, and so also at most \( x^{1 - \frac{1}{k+1} + o(1)} \) possibilities for the \( k \)-set, by Lemma 4.1. So we suppose that \( \text{rad}(B) > x^{1 - \frac{1}{k+1}} \).

Let \( p \) be a prime dividing \( B \). Then \( p \) divides each \( n_i \). Since \( n_1, \ldots, n_k \) is primitive friendly, it is impossible for \( p \) to exactly divide each \( n_i \). Thus, if \( N = n_1 \cdots n_k \), then \( \text{rad}(B)^{k+1} \mid N \). Hence, the number of possibilities for the integer \( N \leq x^k \) is at most

\[ \sum_{r > x^{1 - \frac{1}{k+1}}} \frac{x^k r^{k+1}}{r^{k+1}} \ll x^{1 - \frac{1}{k+1}}. \]

From the maximal order of the \( k \)-fold divisor function, \( N \) determines the set \( \{n_1, \ldots, n_k\} \) in at most \( x^{o(1)} \) ways. The result follows. \hfill \Box

\textbf{Proof of Theorem 1.6.} We may assume \( k \geq 1 \), since when \( k = 0 \) the asymptotic claimed in the theorem holds with \( \alpha_0 = 1 \). We begin by placing each primitive \((k+1)\)-set in increasing order, \( n_0 < n_1 < \cdots < n_k \). We then list the primitive \((k+1)\)-sets in increasing order of \( n_k \), breaking ties arbitrarily. Say the \( j \)th tuple is labeled as \( n_{0,j}, \ldots, n_{k,j} \). Put

\[ \mathcal{M}_{i,j}(x) := \{dn_{i,j} : 1 \leq d \leq x/n_{k,j} \text{ and } \gcd(d, \prod_{t=0}^{k} n_{t,j}) = 1 \}, \]

and observe that a natural number \( m \leq x \) has at least \( k \) friends in \([1, x]\) if and only if \( m \) belongs to one of the sets \( \mathcal{M}_{i,j}(x) \). For each \( J \geq 1 \), put

\[ \alpha_k^{(J)} = \lim_{x \to \infty} \frac{1}{x} \cdot \# \bigcup_{1 \leq j \leq J} \mathcal{M}_{i,j}(x). \]

The existence of the limit follows from inclusion–exclusion. (Each finite intersection of the sets \( \mathcal{M}_{i,j}(x) \) can be described as the set of numbers not exceeding a certain constant.
multiple of \( x \) that satisfy certain congruence conditions.) The \( \alpha^{(J)}_k \) are weakly increasing in \( J \) and bounded by 1; thus, it makes sense to set
\[
\alpha_k = \lim_{J \to \infty} \alpha^{(J)}_k.
\]
Clearly, \( \liminf_{x \to \infty} N_k(x)/x \geq \alpha_k \). We now show that \( \limsup_{x \to \infty} N_k(x)/x \leq \alpha_k \).

If \( m \leq x \) has at least \( k \) friends in \([1, x]\), but \( m \not\in \bigcup_{0 \leq j \leq k} \mathcal{M}_{i,j}(x) \), then \( m \in \mathcal{M}_{i,j}(x) \) for some \( j > J \). Since \( \# \mathcal{M}_{i,j}(x) \leq \frac{x}{n_{k,j}} \), the number of these \( m \leq x \) is at most
\[
(k + 1)x \cdot \sum_{j > J} \frac{1}{n_{k,j}}.
\]
We now study the sum on \( j \). Without the restriction to \( j > J \), the sum converges (or is a finite sum). Indeed, from Lemma 5.1, the contribution to the unrestricted sum from those \( j \) with \( X \leq n_{k,j} < 2X \) is bounded by \( X^{-\frac{1}{k+2} + o(1)} \), as \( X \to \infty \); the desired result now follows upon summing \( X \) along the powers of 2. So with the restriction to \( j > J \), the series decays to 0 as \( J \to \infty \). Letting \( x \to \infty \) and then \( J \to \infty \), in that order, we get that \( \limsup_{x \to \infty} N_k(x)/x \leq \alpha_k \).

It remains to prove that \( \alpha_k \to 0 \). Since \( k \cdot N_k(x) \leq \sum_{n \leq x} F(n; x) \), Erdős’s estimate (2) shows that \( \alpha_k \ll 1/k \).

\[\square\]

Remark.

(i) A similar argument to the one just seen will show that each positive integral moment \( \frac{1}{x} \sum_{n \leq x} F(n; x)^k \) tends to a finite limit as \( x \to \infty \). Here one uses that \( F(n; x)^k \) can be viewed as the count of (ordered) friendly \((k + 1)\)-tuples in \([1, x]\) having final element \( n \) and all other elements distinct from \( n \).

As a consequence, the \( \alpha_k \) tend to zero more rapidly than any power of \( k^{-1} \).

(ii) The computations of \( N_i(x) \), \( i = 1, \ldots, 4 \), reported in Table 3 were made with the aforementioned \texttt{solveBA} script.

We do not know how to show that each \( \alpha_k \) is nonzero. This is equivalent to the assertion that the count of natural numbers \( n \) with \( \frac{\sigma(n)}{n} = \alpha \) is unbounded in \( \alpha \), a conjecture mentioned already in [19]. Note that this conjecture would follow from the infinitude of perfect numbers. Since the \( \alpha_k \) may eventually vanish, we cannot show that the \( \alpha_k \) are strictly decreasing. However, we can prove that the sequence is strictly decreasing until it hits zero.

**Theorem 5.2.** If \( \alpha_k > 0 \), then \( \alpha_k > \alpha_{k+1} \).

**Proof.** Since \( \alpha_k > 0 \), there are examples of friendly \((k + 1)\)-sets. (In the case \( k = 0 \), we consider every singleton subset of the natural numbers to be a friendly \((k + 1)\)-set.) Fix such a set \( \{n_0 < n_1 < \cdots < n_k\} \) once and for all, chosen to have \( n_k \) as small as possible. This set is automatically primitive. Every integer of the form \( m = dn_k \), with \( d \leq x/n_k \) and gcd\((d, n_0 \cdots n_k) = 1\), has at least \( k \) friends in \([1, x]\). It suffices to show that \( \gg x \) values of \( d \) give rise to an \( m \) with exactly \( k \) friends in \([1, x]\).

Sort all of the primitive friendly \((k + 2)\)-sets in increasing order, and then list the sets in order of their largest entry. Say the \( j \)th set consists of \( n_{0,j} < n_{1,j} < \cdots < n_{k+1,j} \). With \( J \) a large, fixed integer to be chosen later, we restrict attention to \( d \) satisfying
\[
(24) \quad \frac{x}{n_k + 1} < d \leq \frac{x}{n_k} \quad \text{and} \quad d \equiv 1 \pmod{\prod_{i=0}^{k} n_i \prod_{1 \leq j \leq J} n_{i,j}}.
\]

Let \( M_J \) be the modulus of this congruence. For large \( x \), there are \( \gg x/M_J \) values of \( d \) satisfying (24). (The implied constant also depends on \( k \) and the set \( \{n_0, \ldots, n_k\} \), but
Hence, the number of integers to consider, namely $dn_k$ for $d < J$, is at most $xM^{-\theta} / \log(nk)$.

In particular, $n_{i,j} | dn_k$. Suppose to begin with that $j \leq J$. Then $\gcd(n_{i,j}, d) = 1$ from (24), and so $n_{i,j} | n_k$. Putting $e = n_k/n_{i,j}$, we have $d' = de \leq x/n_{k+1,j}$, so that $d \leq x/e \cdot n_{k+1,j} \leq x/n_{k+1,j}$.

Since $n_0,j,\ldots,n_{k,j}$ are in increasing order and make up a friendly $(k + 1)$-set, we have $n_{k+1,j} > n_{k,j} \geq n_k$, by the initial choice of $n_k$. Hence, $n_{k+1,j} \geq n_k + 1$, and so this last upper bound on $d$ contradicts (24). So we must have $j > J$. From $dn_k = d'n_{i,j}$ and $\gcd(d, M_j) = 1$, we see that $\gcd(n_{i,j}, M_j) | n_k$. In particular, $\gcd(n_{i,j}, M_j) \ll 1$. Using $dn_k = d'n_{i,j}$ and (24), we see that $d'$ belongs to a fixed residue class modulo $n_k M_j / \gcd(n_{i,j}, n_k M_j)$. The number of these $d' \leq x/n_{k+1,j}$ is at most

$$x \frac{\gcd(n_{i,j}, n_k M_j)}{n_k M_j \cdot n_{k+1,j}} + 1 \ll \frac{x}{M_j \cdot n_{k+1,j}} + 1.$$ 

Hence, the number of $m = dn_k$ with $d$ satisfying (24) that have at least $k + 1$ friends in $[1, x]$ is

$$\ll \sum_{i=0}^{k+1} \sum_{j > J \atop n_{k+1,j} \leq x} \left( \frac{x}{M_j \cdot n_{k+1,j}} + 1 \right).$$

Since the sum on $j$ is only over $j > J$, and since we can take $J$ arbitrarily large, the main term here can be made smaller than any constant multiple of $x/M_j$, by the same argument used in the proof of Theorem 1.6. The error term is bounded by a power of $x$ less than 1, by Lemma 5.1. So fixing $J$ large enough, we have that for large $x$, there are $\gg x/M_j \gg x$ values of $d$ satisfying (24) for which $m = dn_k$ has exactly $k$ friends in $[1, x]$, as desired. 

\[ \Box \]

6. Some final remarks

In the paper [17], which inspired our work above on friendly pairs, Erdős writes “I would like to call attention to three simple problems which as far as I know are still unsolved.”

The problems are:

- Are there infinitely many solutions to the equation $\sigma(n) = \varphi(m)$?
- For each real number $c \geq 1$, is there an infinite sequence of pairs $m_k, n_k$ with $\sigma(m_k) = \sigma(n_k)$ and $n_k/m_k \to c$ as $k \to \infty$?
- If $g(x)$ is the number of coprime pairs $a, b$ with $a < b \leq x$ and $\sigma(a) = \sigma(b)$, do we have $g(x)/x \to \infty$ as $x \to \infty$?

We note that the first two problems have been recently solved in the affirmative (see [24], [33]), and the third was “almost” solved in the affirmative (solved without the coprimality condition) in 1980 (see [36]). It seems fitting to close this paper with a complete proof of the third problem.

Following an idea of Erdős [14], say $\theta > 1$ is such that there are $y^\theta/(\log y)^{O(1)}$ primes $p \leq y^\theta$ with $P(p+1) \leq y$. Let $x$ be large, let $y = \log x$, let $\mathcal{S}$ be the set of primes $p \leq y^\theta$ with $P(p+1) \leq y$, let $n = \#\mathcal{S}$, and let $k = \lfloor \log x/\log(y^\theta) \rfloor$. Consider integers $a$ which are the product of $k$ distinct primes in $\mathcal{S}$. Since $\binom{k}{k} = x^{-1/\theta + o(1)}$, as $x \to \infty$, there are $x^{1-1/\theta + o(1)}$ of these integers $a \leq x$, each with $P(\sigma(a)) \leq y$. But, as was shown by Erdős (see [11]), the number of integers to $x$ not divisible by any prime $> y$ is $x^{o(1)}$. Hence, there
is one value $v$ such that $\sigma(a) = v$ has $x^{1-1/\theta+o(1)}$ solutions $a \leq x$, as $x \to \infty$. It follows that there are $x^{2-2/\theta+o(1)}$ pairs $a < b \leq x$ with $\sigma(a) = \sigma(b)$.

This would solve the third problem above if we had $\theta > 2$ and if the pairs $a, b$ created were coprime. We know from [36] that $\theta > 2$, the current record being a number slightly larger than 3.377, see [4]. Let $\epsilon > 0$ be fixed but arbitrarily small, and let $k_0 = \lfloor (1-\epsilon)k \rfloor$. The number of pairs $a, b$ which reduce to a coprime pair $a', b'$ with $\omega(a') = \omega(b') = j$ is at most

$$\binom{n}{j} \binom{n}{k-j}.$$ 

Summing this for $j \leq k_0$ gives $x^{2-2/\theta-(1-1/\theta)+o(1)}$. Thus, almost all of our pairs $a, b$ reduce to a pair $a', b'$ where $\omega(a') = \omega(b') > k_0$. In fact, there is some $j_0 \in (k_0, k]$ where $x^{2-2/\theta+o(1)}$ pairs $a, b$ reduce to coprime pairs $a', b'$ with $\omega(a') = \omega(b') = j_0$. Each of these pairs comes from at most $\binom{n}{k-j_0} \leq x^{(1-1/\theta)+o(1)}$ pairs $a, b$, so there must be at least $x^{(2-\epsilon)(1-1/\theta)+o(1)}$ coprime pairs $a', b' \leq x$ with $\sigma(a') = \sigma(b')$. By choosing $\epsilon$ sufficiently small, we see that $g(x) > x^{1.4}$ for all sufficiently large values of $x$.

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