What we still don't know about addition and multiplication

Carl Pomerance, Dartmouth College Hanover, New Hampshire, USA You would think that all of the issues surrounding addition and multiplication were sewed up in third grade!

Well in this talk we'll learn about some things they didn't tell you ...

Here's one thing they did tell you:

Find 483×784 .

× 784

If instead you had a problem with two 23-digit numbers, well you always knew deep down that math teachers are cruel and sadistic. Just kidding! (*Aside: evil laugh* ...)

In principle if you really have to, you could work out 23-digits times 23-digits on paper, provided the paper is big enough, but it's a lot of work.

So here's the real question: How much work?

Of course the amount of work depends not only on how long the numbers are, but on what they are. For example, multiplying 10^{22} by 10^{22} , that's 23-digits times 23-digits, but you can do it in your head.

In general, you'll take each digit of the lower number, and multiply it painstakingly into the top number. It's less work if some digit in the lower number is repeated, and there are definitely repeats, since there are only 10 possible digits. But even if it's no work at all, you still have to write it down, and that's 23 or 24 digits. At the minimum (assuming no zeroes), you have to write down $23^2 = 529$ digits for the "parallelogram" part of the product. And then comes the final addition, where all of those 529 digits need to be processed.

So in general if you multiply two n-digit numbers, it would seem that you'd be taking n^2 steps, unless there were a lot of zeroes. This ignores extra steps, like carrying and so on, but that at worst changes n^2 to maybe $2n^2$ or $3n^2$. We say that the "complexity" of "school multiplication" for two n-digit numbers is of order n^2 .

Here is what we don't know:

What is the fastest way to multiply?

There's a method known as the Fast Fourier Transform that allows you to multiply in a lot fewer than n^2 steps, but more than any fixed constant times n steps. We don't know if the Fast Fourier Transform is the best possible.

Let's play Jeopardy Multiplication!

Here are the rules: I give you the answer to the multiplication problem, and you give me the problem phrased as a question. You must use whole numbers larger than 1.

So, if I say "15", you say "What is 3×5 ?"

OK, let's play.

21

Good. That was easy. Let's up the ante.

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91

What is 7×13 ?

Let's do 8051.

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Thinking, thinking Hmm,

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$$8051 = 8100 - 49 = 90^2 - 7^2 = (90 - 7)(90 + 7) = 83 \times 97.$$

Got it!

What is 83×97 ?

So, here's what we don't know:

How many steps does it take to figure out the factors if you are given an n-digit number which can be factored? (A trick problem would be: 17. The only way to write it as $a \times b$ is to use 1, and that was ruled out. So, prime numbers cannot be factored, and the thing we don't know is how long it takes to factor the non-primes.)

The best answer we have so far is about $10^{n^{1/3}}$ steps, and even this is not a theorem, but our algorithm (the *number field sieve*) seems to work in practice.

This is all crucially important for the security of Internet commerce. Or I should say that Internet commerce relies on the premise that we *cannot* factor much more quickly than that.

Lets look at an unsolved problem concerning addition.

We all recall the addition table:

+	1	2	3	4	5	6	7	8	9	10
1	2	3	4	5	6	7	8	9	10	11
2	3	4	5	6	7	8	9	10	11	12
3	4	5	6	7	8	9	10	11	12	13
4	5	6	7	8	9	10	11	12	13	14
5	6	7	8	9	10	11	12	13	14	15
6	7	8	9	10	11	12	13	14	15	16
7	8	9	10	11	12	13	14	15	16	17
8	9	10	11	12	13	14	15	16	17	18
9	10	11	12	13	14	15	16	17	18	19
10	11	12	13	14	15	16	17	18	19	20

The 10×10 array of sums has all the numbers from 2 to 20 for a total of 19 different sums.

If you were to try this for the $N \times N$ addition table we'd see all of the numbers from 2 to 2N for a total of 2N-1 different sums.

Now, what if we were to be perverse and instead of having the numbers from 1 to N, we had some arbitrary list of N different numbers.

What is the *fewest* number of different sums that you can make?

After a little reflection you answer 2N-1, and you'd be right.

Now consider the reverse problem. You have a list of N different numbers, you form the addition table with themselves, and you find that there are just 2N-1 different sums.

Must the list be the numbers from 1 to N?

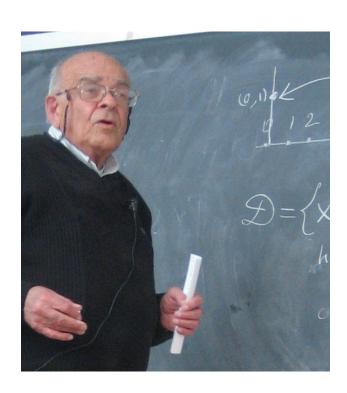
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×	2	5	8	11	14
2	4	7	10	13	18
5	7	10	13	16	19
8	10	13	16	19	22
11	13	16	19	22	25
14	16	19	22	25	28

There are the 9 sums: 4, 7, 10, 13, 16, 19, 22, 25, 28

Gregory Freiman: If there are just 2N-1 distinct sums, then the N-element set must be an arithmetic progression. More generally, if there are at most cN distinct sums, where c is fixed, then the set is contained within some low-dimensional generalized arithmetic progression.



Let L(N) denote the number of (decimal) digits of N.

What we still don't know:

What structure is imposed on a list of N numbers if there are only at most NL(N) distinct pairwise sums?

What about 3-way sums?

Here's something with multiplication tables.

Let's look at the $N \times N$ multiplication table using the numbers from 1 to N. With addition we were able to count exactly how many distinct numbers appear in the table.

How many different numbers appear in the $N \times N$ multiplication table?

Let M(N) be the number of distinct entries in the $N\times N$ multiplication table.

×	1	2	3	4	5
1	1	2	3	4	5
2	2	4	6	8	10
3	3	6	9	12	15
4	4	8	12	16	20
5	5	10	15	20	25

So, M(5) = 14.

M(10) = 42.

It may be too difficult to expect a neat exact formula for M(N).

Instead, we could ask for its order of magnitude, or even approximate order of magnitude.

For example, does M(N) go to infinity like a constant times N^2 , or more slowly. That is, maybe there is a positive number c with

$$M(N) > cN^2$$

for all N. Or maybe for *every* positive number c,

$$M(N) < cN^2$$

for infinitely many choices for N or even for all large N.

Here are some values of $M(N)/N^2$ (Brent & Kung 1981):

N	M(N)	$M(N)/N^2$
1	1	1.0000
3	6	0.6667
7	25	0.5102
15	89	0.3956
31	339	0.3528
63	1237	0.3117
127	4646	0.2881
255	17577	0.2703
511	67591	0.2588
1023	258767	0.2473
2047	1004347	0.2397
4095	3902356	0.2327
8191	15202049	0.2266

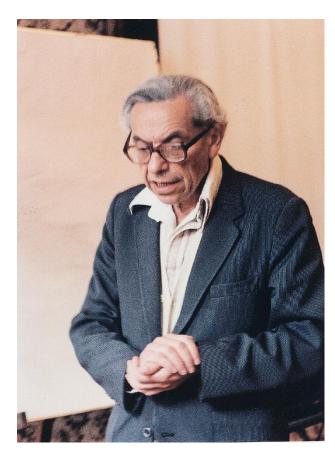
And some more values (Brent & Kung 1981, Brent 2012):

N	M(N)	$M(N)/N^2$
$2^{14} - 1$	59410556	0.2213
$2^{15} - 1$	232483839	0.2165
$2^{16} - 1$	911689011	0.2123
$2^{17} - 1$	3581049039	0.2084
$2^{18} - 1$	14081089287	0.2049
$2^{19} - 1$	55439171530	0.2017
$2^{20} - 1$	218457593222	0.1987
$2^{21} - 1$	861617935050	0.1959
$2^{22}-1$	3400917861267	0.1933
$2^{23} - 1$	13433148229638	0.1909
$2^{24} - 1$	53092686926154	0.1886
$2^{25} - 1$	209962593513291	0.1865

And some statistically sampled values (Brent & P 2012):

	$(N)/N^2$
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	0.0348 0.0312 0.0269 0.0240 0.0216 0.0186 0.0171 0.0153 0.0133 0.0122 0.0115 0.0095

Paul Erdős studied this problem in two papers, one in 1955, the other in 1960.



Paul Erdős, 1913–1996

In 1955, Erdős proved (in Hebrew) that $M(N)/N^2$ tends to 0 as N runs off to infinity.

In 1960, at the prodding of Linnik and Vinogradov, Erdős identified (in Russian) how quickly $M(N)/N^2$ tends to 0. Recall that L(N) denotes the number of (decimal) digits of N. There is a constant (let's call it E for Erdős),

$$E = 0.08607133...,$$

such that M(N) is about $N^2/L(N)^E$.

What's with this weasel word "about"? We're saying that M(N) as a first approximation is given by $N^2/L(N)^E$, and no other exponent on L(N) does better. We're not saying that dividing one by the other, it must tend to 1.

In work of Tenenbaum progress was made (in French) in nailing this down.

In 2008, Ford showed (in English) that M(N) is of order of magnitude

$$\frac{N^2}{L(N)^E L(L(N))^{3/2}}.$$

No matter the language,

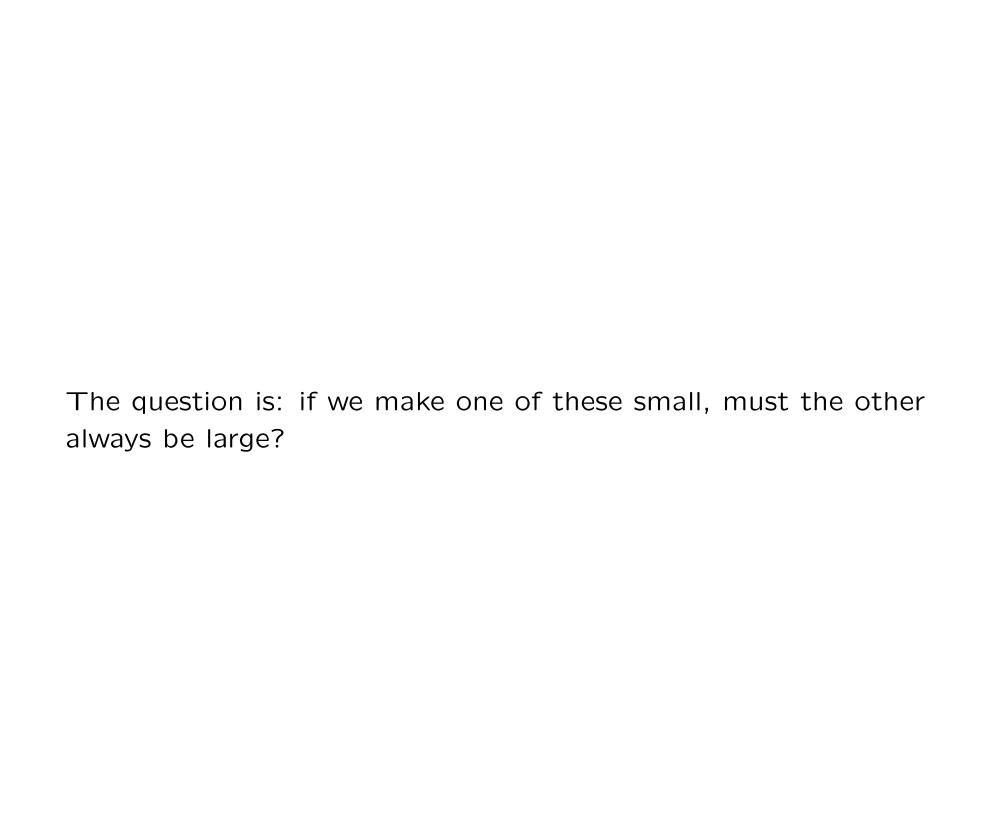
we still don't know an asymptotic estimate for M(N), despite this just being about multiplication tables!

Let me close with one unified problem about addition and multiplication together. It's due to Erdős & Szemerédi.

Look at **both** the addition and multiplication tables and count the number of distinct entries.

We've seen that if we take the first N numbers we get about $N^2/L(N)^E$ entries (since the 2N-1 from the addition table is just "noise" compared to the huge number from the multiplication table).

At the other extreme, if we take for our N numbers the powers of 2, namely $1, 2, 4, \ldots, 2^{N-1}$, then there are at least $\frac{1}{2}N^2$ distinct entries in the addition table and only 2N-1 entries in the multiplication table.



Put more precisely: If we have N distinct numbers must the number of distinct pairwise sums plus the number of distinct pairwise products be greater than $N^{1.999}$ for all large values of N?

Or just ask for more than $N^{1.5}$?

We don't know.

And it's not for lack of trying.

The game players with the sum/product problem: Erdős, Szemerédi, Nathanson, Chen, Elekes, Bourgain, Chang, Konyagin, Green, Tao, Solymosi, ...

The best that's been proved (Solymosi) is that there are at least $N^{4/3}$ different entries.

This list of mathematicians contains two Fields Medalists, a Wolf Prize winner, an Abel Prize winner, four Salem Prize Winners, two Crafoord Prize winners, and an Aisenstadt Prize winner.

And still the problem is not solved!

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THANK YOU