

SQUARE VALUES OF EULER'S FUNCTION

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ABSTRACT. We show that almost all squares are missing from the range of Euler's φ -function.

1. INTRODUCTION

Let φ denote Euler's function, let \mathbf{N} denote the set of positive integers, and let $\mathcal{V} = \varphi(\mathbf{N})$, the set of values of φ . Further, let $V(x) = \#\{n \leq x : n \in \mathcal{V}\}$. The distribution of \mathcal{V} has been of interest since the 1930s when Erdős showed that $V(x) = x/(\log x)^{1+o(1)}$ as $x \rightarrow \infty$. We still do not have an asymptotic for $V(x)$, but after work of Ford [8], we do know the order of magnitude.

For a function $f: \mathbf{N} \rightarrow \mathbf{N}$, let

$$\mathcal{V}_f = \{n : f(n) \in \mathcal{V}\}, \quad \mathcal{V}^f = \{n : \varphi(n) \in f(\mathbf{N})\},$$

and let $V_f(x), V^f(x)$ be the respective counting functions for $\mathcal{V}_f, \mathcal{V}^f$. The situation when f is a linear polynomial is fairly well-understood. If $f(n) = kn$, where k is a fixed natural number, then $V_f(x) \sim V(kx)$ and $V^f(x) \sim x$ as $x \rightarrow \infty$; on the other hand, if $f(n) = kn + j$ with $0 < j < k$, then $V_f(x) = o(V(kx))$ and $V^f(x) = o(x)$. (The V_f -results do not appear to be in the literature, but follow from the method of Ford.) More refined results concerning the cases when $0 < j < k$ can be found in [16, 7, 9]. The case when $f = \sigma$, the sum-of-divisors function, was considered in [10], where some old questions of Erdős were settled (see also [11, 12]). This paper is concerned with the function $f(n) = n^2$, which we denote with the symbol \square , so that

$$V_{\square}(x) = \#\{n \leq x : n^2 \in \mathcal{V}\}, \quad V^{\square}(x) = \#\{n \leq x : \varphi(n) = m^2 \text{ for some integer } m\}.$$

It was shown in [2], perhaps counter-intuitively, that $V^{\square}(x) \geq x^{0.7}$ for all large x , with the conjectured exponent on x allowed to be any number below 1. In that paper it was also shown that $V_{\square}(x) \geq x^{0.234}$ for all sufficiently large x . This lower bound was considerably improved in [3], where it was shown that $V_{\square}(x) \gg x/(\log x)^4$ (compare with the case $r = 2$ of [13, Theorem 1.2]).

The paper [2] shows that $V^{\square}(x) \leq x/\exp((1 + o(1))(\log x \log \log x)^{1/2})$ as $x \rightarrow \infty$, but does not address an upper bound for $V_{\square}(x)$. It is not immediately clear that $V_{\square}(x) = o(x)$. In fact, a short computer run shows that $V_{\square}(10^8) = 26,094,797$ so that more than half of the even numbers to 10^8 have their squares in the range of φ . In this paper we prove the following results.

Theorem 1. *For all sufficiently large numbers x , we have $V_{\square}(x) \leq x/(\log x)^{.0063}$.*

Theorem 2. *We have $V_{\square}(x) \gg x/(\log x)^3$.*

In addition, we discuss some heuristics for the estimation of $V_{\square}(x)$ and we discuss the analogous problems for the sum-of-divisors function.

It would be interesting to obtain versions of Theorems 1 and 2 with φ replaced by the Carmichael λ -function, but we have not so far succeeded in this.

Notation. We use the Landau/Bachmann O and o -notation, as well as the associated Vinogradov \ll and \gg notations, with their standard meanings. We write $A \asymp B$ to mean that $A \ll B$ and $B \ll A$. Any dependence of implied constants is noted explicitly, often with a subscript.

The letters p , q , and ℓ , with or without subscripts, always denote primes. We use $P(n)$ for the largest prime factor of the natural number n , with the convention that $P(1) = 1$. The notation $p^e \parallel n$ means that $p^e \mid n$ but that $p^{e+1} \nmid n$; in this case, we say that p^e *exactly divides* n . As usual, $\Omega(n)$ denotes the number of prime factors of n counted with multiplicity; thus, $\Omega(n) = \sum_{p^k \parallel n} k$.

We write \log_k for the k -fold iterate of the natural logarithm.

2. PREPARATION

2.1. Anatomy and sieving. A classical theorem of Hardy and Ramanujan asserts that a typical natural number n has about $\log_2 n$ prime factors, regardless of whether or not the primes are counted with multiplicity. Our first lemma, which may be deduced from the results in Chapter 0 of [15], bounds from above the number of n for which $\Omega(n)$ is atypically large.

Lemma 3. *Let $x \geq 3$, and let $\epsilon > 0$. For $1 \leq \alpha \leq 2 - \epsilon$, the number of $n \leq x$ with $\Omega(n) \geq \alpha \log_2 x$ is $O_\epsilon(x(\log x)^{-Q(\alpha)})$, where we set $Q(\lambda) = \int_1^\lambda \log t \, dt = \lambda \log(\lambda) - \lambda + 1$.*

We now quote two upper bound sieve results, in slightly crude forms that are convenient for our later applications. Both of these follow from the general upper bound O -result appearing as [14, Theorem 2.2].

Lemma 4. *Suppose that A_1, \dots, A_h are positive integers and B_1, \dots, B_h are integers such that*

$$E := \prod_{i=1}^h A_i \prod_{1 \leq i < j \leq h} (A_i B_j - A_j B_i) \neq 0.$$

Then for $x \geq 3$,

$$\#\{n \leq x : A_i n + B_i \text{ prime for all } 1 \leq i \leq h\} \ll \frac{x}{(\log x)^h} (\log_2 |3E|)^h,$$

where $\nu(p)$ is the number of solutions of the congruence $\prod (A_i n + B_i) \equiv 0 \pmod{p}$, and the implied constant may depend on h .

Lemma 5. *Let A , B , and C be integers with $A > 0$ and $D = B^2 - 4AC$ not a square. Write $D = df^2$, where d is a fundamental discriminant. Then for $x \geq 3$,*

$$\#\{p \leq x : Ap^2 + Bp + C \text{ prime}\} \ll \frac{x}{(\log x)^2} (\log_2 |3ACD|)^3 \prod_{\ell \leq x} \left(1 - \frac{\left(\frac{d}{\ell}\right)}{\ell}\right), \quad (1)$$

where $\left(\frac{d}{\cdot}\right)$ is the Kronecker symbol.

2.2. Sieving quadratics and short Euler products. To control the size of the product on ℓ appearing in (1), we appeal to the methods and results of a recent preprint of Chandee, David, Koukoulopoulos, and Smith [5].

Lemma 6. *Let $\epsilon > 0$. Let χ be a nonprincipal real character mod q . For all real $y \geq 1$, we have*

$$\prod_{\ell \leq y} \left(1 - \frac{\chi(\ell)}{\ell}\right) \ll_\epsilon q^\epsilon.$$

Proof. The proof parallels that of [5, Lemma 3.2]. By Mertens' theorem, $\prod_{\ell \leq y} (1 - \chi(\ell)/\ell) \ll q^\epsilon \prod_{\exp(q^\epsilon) < \ell \leq y} (1 - \chi(\ell)/\ell)$. Hence, it suffices to show that the remaining product is $O_\epsilon(1)$. By the classical Siegel–Walfisz estimates (see [6, eq. (3), p. 132]),

$$\sum_{n \leq x} \Lambda(n) \chi(n) \ll_\epsilon x / \log x \quad \text{for all } x \geq \exp(q^\epsilon). \quad (2)$$

Recalling that $\log(1 - t) = -\sum_{k \geq 1} t^k/k$ (for $|t| < 1$), we find that

$$\begin{aligned} \log \prod_{\exp(q^\epsilon) < \ell \leq y} \left(1 - \frac{\chi(\ell)}{\ell}\right) &= - \sum_{n > 1, \ell | n \Rightarrow \exp(q^\epsilon) < \ell \leq y} \frac{\Lambda(n) \chi(n)}{n \log n} \\ &= - \sum_{\exp(q^\epsilon) < n \leq y} \frac{\Lambda(n) \chi(n)}{n \log n} + O(1) \ll_\epsilon 1. \end{aligned}$$

where the final estimate is obtained from (2) by partial summation. \square

The next lemma is an equivalent form of [5, Lemma 3.3], which the authors of that paper attribute in essence to Elliott.

Lemma 7. *Fix $\delta \in (0, 1]$, and let $Q \geq 3$. We can choose a set $\mathcal{E}_\delta(Q)$ of real, primitive characters, all of conductor bounded by Q , with*

$$\#\mathcal{E}_\delta(Q) \ll_\delta Q^\delta$$

and so that the following holds: If χ is a primitive real character of conductor $q \leq Q$ and $\chi \notin \mathcal{E}_\delta(Q)$, then

$$\prod_{y < \ell \leq z} \left(1 - \frac{\chi(\ell)}{\ell}\right) \asymp_\delta 1 \quad \text{uniformly for } z \geq y \geq \log Q.$$

For each nonsquare integer d , let χ_d be the primitive real character of conductor $|D|$ given by the Kronecker symbol $\left(\frac{D}{\cdot}\right)$, where D is the discriminant of $\mathbf{Q}(\sqrt{d})$. It is convenient for us to isolate the following consequence of Lemma 7.

Lemma 8. *Let \mathcal{D} be the set of squarefree $d \neq 1$ for which there exists a real number y with*

$$\prod_{\ell \leq y} \left(1 - \frac{\chi_d(\ell)}{\ell}\right) \geq (\log_2 |3d|)^2. \quad (3)$$

For fixed $\delta > 0$ and all $x \geq 1$, we have that

$$\#\{d \in \mathcal{D} : |d| \leq x\} \ll_\delta x^\delta.$$

Proof. We can assume that x is large. It suffices to prove the stated estimate for $\#\{d \in \mathcal{D} : x^\delta < |d| \leq x\}$. Let $\{y_i\}_{i=0}^\infty$ be the sequence of real numbers defined by $y_i = 4^i x^\delta$, and choose j so that $y_j < |d| \leq y_{j+1}$. Then the conductor of χ_d is bounded by $4y_{j+1}$, and $4y_{j+1} < 16|d|$. We claim that if $\chi_d \notin \mathcal{E}_\delta(4y_{j+1})$, then the inequality (3) never holds. Indeed, Lemma 7 (with $Q := 4y_{j+1}$) shows that for every y ,

$$\prod_{\ell \leq y} \left(1 - \frac{\chi_d(\ell)}{\ell}\right) \ll_\delta \prod_{\ell \leq \min\{\log(4y_{j+1}), y\}} \left(1 - \frac{\chi_d(\ell)}{\ell}\right) \ll \log_2 |d|,$$

using Mertens' theorem in the final step. Since $|d| \geq x^\delta$ and x is large, this upper bound is incompatible with (3), proving our claim. Since distinct squarefree d give rise to distinct

primitive real characters χ_d , the upper bound for $\#\mathcal{E}_\delta(Q)$ from Lemma 7 yields

$$\#\{d \in \mathcal{D} : x^\delta < |d| \leq x\} \leq \sum_{\substack{j \geq 0 \\ y_j \leq x}} \#\mathcal{E}_\delta(4y_{j+1}) \ll_\delta \sum_{0 \leq j \leq \frac{\log(x^{1-\delta})}{\log 4}} 4^{(j+2)\delta} x^{\delta^2} \ll_\delta x^\delta.$$

This completes the proof of the lemma. \square

3. PROOF OF THE UPPER BOUND (THEOREM 1)

Setup. We assume throughout the argument that x is large. Let $n \leq x$ be such that $n^2 = \varphi(m)$ for some integer m . By de Bruijn [4, eq. (1.6)], we can assume that

$$(i) \ P(n) \geq x^{1/\log_2 x}$$

since the number of $n \leq x$ for which (i) fails is $O(x/\log x)$. We can also assume that

$$(ii) \ n \text{ is not divisible by any } d \in \mathcal{D} \text{ with } |d| > \log x, \text{ where } \mathcal{D} \text{ is the set considered in Lemma 8.}$$

Indeed, since $\#\{d \in \mathcal{D} : |d| \leq t\} \ll t^{1/2}$ for all $t \geq 1$, the count of exceptional $n \leq x$ is $O(x/(\log x)^{1/2})$ (by partial summation). At the cost of an additional exceptional set of the same order, we can further assume that

$$(iii) \ n \text{ is not divisible by any square exceeding } \log x.$$

Introducing another exceptional set of size $O(x/(\log x)^{1/2})$, we can assume that

$$(iv) \ \text{there is no prime } p^2 \text{ dividing } m \text{ with } p > \log x.$$

Indeed, suppose that $p^2 \mid m$. Setting $r_p = \prod_{\ell^e \parallel p-1} \ell^{\lfloor e/2 \rfloor}$, we see that $p \cdot r_p \mid n$. Note that $r_p \geq \sqrt{p-1} \gg \sqrt{p}$. Hence, the number of n with $p^2 \mid m$ for some $p > \log x$ does not exceed

$$\sum_{p > \log x} \frac{x}{p \cdot r_p} \ll x \sum_{p > \log x} \frac{1}{p^{3/2}} \ll x/(\log x)^{1/2}.$$

Let α be a parameter with $1 < \alpha < 2$, which will be chosen later so as to optimize the argument. We may assume that

$$(v) \ \Omega(n) \leq \alpha \log_2 x,$$

since Lemma 3 guarantees that the number of exceptions $n \leq x$ is

$$\ll_\alpha x/(\log x)^{1-\alpha+\alpha \log \alpha}. \quad (4)$$

Let $p = P(n)$, so that $p^2 \mid n^2 = \varphi(m)$. By (i) and (iv), we have that $p^2 \nmid m$, and so there are only two ways to explain how $p^2 \mid \varphi(m)$:

- I. there are two different primes $q_1, q_2 \mid m$ with $q_i \equiv 1 \pmod{p}$ for $i = 1, 2$,
- II. there is a prime $q \mid m$ with $q \equiv 1 \pmod{p^2}$.

Case I. We will assume that the primes q_1, q_2 are not $1 \pmod{p^2}$; otherwise we may push this situation into Case II. For such a prime q we may write it as $1 + apb^2$, where ap is squarefree. This shows that n may be written in the form

$$n = ua_1a_2a_3b_1b_2p, \quad \text{with } a_1a_2a_3p \text{ squarefree, } 1 + a_1a_3pb_1^2 \text{ prime, } 1 + a_2a_3pb_2^2 \text{ prime.}$$

For each fixed choice of $u, a_1, a_2, a_3, b_1, b_2$ we count primes $p \leq x/ua_1a_2a_3b_1b_2$ with the two primality conditions above holding. Using the upper bound sieve in the form of Lemma 4, and recalling that $x/ua_1a_2a_3b_1b_2 \geq p > x^{1/\log_2 x}$, we find that the number of these p is

$$\ll \frac{x}{ua_1a_2a_3b_1b_2(\log x)^3} (\log_2 x)^6. \quad (5)$$

(Explicitly, we apply Lemma 4 with $A_1 = 1$ and $B_1 = 0$, $A_2 = a_1a_3b_1^2$ and $B_2 = 1$, and $A_3 = a_2a_3b_2^2$ and $B_3 = 1$; note that since $q_1 \neq q_2$, we have $E \neq 0$, and $|E| < x^{O(1)}$.) Now we

sum our upper bound (5) over the possibilities for $u, a_1, a_2, a_3, b_1, b_2$, keeping in mind that their product is bounded by x and $\Omega(ua_1a_2a_3b_1b_2) \leq \alpha \log_2 x$. Here it is helpful to introduce an auxiliary parameter z (Rankin's trick); notice that when $0 < z < 1$,

$$\sum \frac{1}{ua_1a_2a_3b_1b_2} \leq z^{-\alpha \log_2 x} \sum \frac{z^{\Omega(u)} z^{\Omega(a_1)} z^{\Omega(a_2)} z^{\Omega(a_3)} z^{\Omega(b_1)} z^{\Omega(b_2)}}{ua_1a_2a_3b_1b_2}.$$

Keeping only the restriction that $P(ua_1a_2a_3b_1b_2) \leq x$, we find that

$$\sum \frac{z^{\Omega(u)} z^{\Omega(a_1)} z^{\Omega(a_2)} z^{\Omega(a_3)} z^{\Omega(b_1)} z^{\Omega(b_2)}}{ua_1a_2a_3b_1b_2} \leq \left(\prod_{\ell \leq x} (1 - z/\ell)^{-1} \right)^6 \ll (\log x)^{6z}.$$

(The last estimate uses a weak form of Mertens' theorem.) Comparing the previous two displays, we find that $\sum \frac{1}{ua_1a_2a_3b_1b_2} \ll (\log x)^{6z - \alpha \log_2 z}$. To optimize, we take $z = \alpha/6$ to get an upper bound of $O((\log x)^{\alpha - \alpha \log(\alpha/6)})$ for our reciprocal sum. Referring back to (5), we see that the total count of n in Case I is

$$\ll \frac{x}{(\log x)^{3 - \alpha + \alpha \log(\alpha/6)}} (\log_2 x)^6. \quad (6)$$

Case II. Write $q - 1 = a(bp)^2$ where a is squarefree, so that $n = uabp$ for some integer u . We first consider the sub-case where $P(ua) \leq \exp((\log x)^\beta)$, where $0 < \beta < 1$ is to be chosen later. For given values of u, a, b , the number of choices for $p \leq x/uab$ satisfying the primality condition is

$$\ll \frac{x}{uab(\log x)^2} (\log_2 x)^5 \prod_{\ell \leq x/uab} \left(1 - \frac{\chi_{-a}(\ell)}{\ell} \right). \quad (7)$$

(Here we have applied Lemma 5 with $A = ab^2$, $B = 0$, and $C = 1$, so that $D = -4ab^2$ and d is the discriminant of $\mathbf{Q}(\sqrt{-a})$.) If $-a \notin \mathcal{D}$, then the product appearing in (7) is $O((\log_2 x)^2)$. If $-a \in \mathcal{D}$, our assumption (ii) implies that $a \leq \log x$. In that case, Lemma 6 shows that the product in (7) is $O_\epsilon((\log x)^{\epsilon/2})$, for any $\epsilon > 0$. So whether or not $-a \in \mathcal{D}$, the number of choices for p is

$$\ll_\epsilon \frac{x}{uab(\log x)^{2 - \epsilon}}. \quad (8)$$

(We have absorbed the powers of $\log_2 x$ into the exponent of $\log x$.) We now sum over u, a, b by the method used in Case I, keeping in mind that $P(ua) \leq \exp((\log x)^\beta)$. For $0 < z < 1$,

$$\begin{aligned} \sum \frac{1}{uab} &\leq z^{-\alpha \log_2 x} \sum \frac{z^{\Omega(u)} z^{\Omega(a)} z^{\Omega(b)}}{uab} \\ &\leq z^{-\alpha \log_2 x} \prod_{\ell_1 \leq x} (1 - z/\ell_1)^{-1} \left(\prod_{\ell_2 \leq \exp((\log x)^\beta)} (1 - z/\ell_2)^{-1} \right)^2 \ll (\log x)^{-\alpha \log_2 z + (1+2\beta)z}. \end{aligned}$$

The optimal choice is $z = \alpha/(1+2\beta)$, which gives $\sum \frac{1}{uab} \ll (\log x)^{\alpha - \alpha \log(\alpha/(1+2\beta))}$. So by (8), the total contribution in this sub-case is

$$\ll_\epsilon \frac{x}{(\log x)^{2 - \alpha + \alpha \log(\alpha/(1+2\beta)) - \epsilon}}. \quad (9)$$

We divide the remaining sub-case when $P(ua) > \exp((\log x)^\beta)$ into further sub-cases as follows. For each positive integer i , let $\beta_i = \beta + i/\log_2 x$, and let \mathcal{I}_i be the interval

$$\mathcal{I}_i = (\exp((\log x)^{\beta_{i-1}}), \exp((\log x)^{\beta_i})].$$

For each i we consider the sub-case where $p_2 := P(ua) \in \mathcal{I}_i$. Clearly, the number of possible sub-cases is at most $1 + \log_2 x$.

We know that $p_2 \mid ua \mid n$, while (iii) implies that $p_2^2 \nmid n$. Hence, $p_2 \parallel n$. Consequently, $p_2 \nmid bp$ and so $p_2^2 \nmid q - 1$. Since $p_2 > \log x$, (iv) gives that $p_2^2 \nmid m$. In conjunction with the relations $p_2^2 \parallel n^2 = \varphi(m)$ and $p_2^2 \nmid q - 1$, this shows that there is a prime $q_2 \neq q$ dividing m with $q_2 \equiv 1 \pmod{p_2}$. If $p_2 \mid u$, then either $p_2^2 \parallel q_2 - 1$ or $p_2 \parallel q_2 - 1$ and there is some other prime $q_3 \mid m$ with $p_2 \parallel q_3 - 1$. If $p_2 \mid a$, then $p_2 \parallel q_2 - 1$. We shall sum up these possibilities as $p_2^k \parallel q_1 - 1$, $k = 0$ or 1 , and $p_2^j \parallel q_2 - 1$, $j = 1$ or 2 and $k + j \leq 2$, ignoring the possible existence of a prime q_3 .

Set $q_1 = q$, $p_1 = p$, $b_1 = b$. We can select natural numbers a_1, a_2, a_3, b_2 with $a_1 a_2 a_3 p_1 p_2$ squarefree and

$$q_1 - 1 = a_1 a_3 b_1^2 p_1^2 p_2^k, \quad q_2 - 1 = a_2 a_3 b_2^2 p_2^j.$$

Then n has a decomposition of the form

$$n = u_1 a_1 a_2 a_3 b_1 b_2 p_1 p_2.$$

Here, in our old notation, $a = a_1 a_3 p_2^k$ and $u = u_1 a_2 b_2 p_2^{1-k}$. Thus, $P(u_1 a_1 a_2 a_3 b_2) < p_2$. Fixing $u_1, a_1, a_2, a_3, b_1, b_2, p_2$ and using the primality of q_1 , we deduce from Lemma 5 (applied with $A = a_1 a_3 b_1^2 p_2^k$, $B = 0$, and $C = 1$) that the number of possible $p_1 \leq x/u_1 a_1 a_2 a_3 b_1 b_2 p_2$ is

$$\begin{aligned} &\ll \frac{x}{u_1 a_1 a_2 a_3 b_1 b_2 p_2 (\log x)^2} (\log_2 x)^5 \prod_{\ell \leq x/u_1 a_1 a_2 a_3 b_1 b_2 p_2} \left(1 - \frac{\chi_{-a_1 a_3 p_2}(\ell)}{\ell}\right) \\ &\ll_{\epsilon} \frac{x}{u_1 a_1 a_2 a_3 b_1 b_2 p_2 (\log x)^{2-\epsilon}}. \end{aligned} \quad (10)$$

(To estimate the product we use an analysis similar to that in (7).) We now fix $u_1, a_1, a_2, a_3, b_1, b_2, p_2$ and sum on $p_2 \in \mathcal{I}_i$. First assume that $j = 1$. Since p_2 and $a_2 a_3 b_2^2 p_2 + 1$ are both prime, the sieve in the form of Lemma 4 shows that for each $t \geq 3$, the number of possible $p_2 \leq t$ is $O(t(\log_2 x)^2/(\log t)^2)$. Now partial summation implies that if we sum (10) over $p_2 \in \mathcal{I}_i$, the result is

$$\ll_{\epsilon} \frac{x}{u_1 a_1 a_2 a_3 b_1 b_2 (\log x)^{2+\beta_{i-1}-2\epsilon}}. \quad (11)$$

(Indeed, this upper bound holds for the larger sum over all $p_2 \geq \exp((\log x)^{\beta_{i-1}})$.) Now assume $j = 2$. We proceed in the same way, though now we use Lemma 5 and a similar analysis as in (7), getting an estimate of

$$\ll_{\epsilon} \frac{x}{u_1 a_1 a_2 a_3 b_1 b_2 (\log x)^{2+\beta_{i-1}-3\epsilon}}. \quad (12)$$

Finally, we replace the estimate (11) with the larger bound (12) and sum over $u_1, a_1, a_2, a_3, b_1, b_2$, keeping in mind that $P(u_1 a_1 a_2 a_3 b_2) \leq \exp((\log x)^{\beta_i})$. For $0 < z < 1$,

$$\begin{aligned} \sum \frac{1}{u_1 a_1 a_2 a_3 b_1 b_2} &\leq z^{-\alpha \log_2 x} \sum \frac{z^{\Omega(u_1)} z^{\Omega(a_1)} z^{\Omega(a_2)} z^{\Omega(a_3)} z^{\Omega(b_1)} z^{\Omega(b_2)}}{u_1 a_1 a_2 a_3 b_1 b_2} \\ &\leq z^{-\alpha \log_2 x} \left(\prod_{\ell_1 \leq \exp((\log x)^{\beta_i})} (1 - z/\ell_1)^{-1} \right)^5 \prod_{\ell_2 \leq x} (1 - z/\ell_2)^{-1} \\ &\ll (\log x)^{-\alpha \log z + (1+5\beta_i)z}. \end{aligned}$$

We select $z = \alpha/(1 + 5\beta_i)$ and find that $\sum \frac{1}{u_1 a_1 a_2 a_3 b_1 b_2} \ll (\log x)^{\alpha - \alpha \log(\alpha/(1+5\beta_i))}$. Referring back to (12), we deduce that the contribution of the i th sub-case is

$$\ll_{\epsilon} \frac{x}{(\log x)^{2+\beta_{i-1}-\alpha+\alpha \log(\alpha/(1+5\beta_i))}-3\epsilon}. \quad (13)$$

To continue our analysis, we make the additional assumption that our parameters α and β satisfy

$$0 < \beta \leq \alpha - \frac{1}{5} \leq 1. \tag{14}$$

As $\beta_i - \beta_{i-1} = 1/\log_2 x$, it is straightforward to check that the upper bound in (13) remains valid with the occurrence of β_i replaced by β_{i-1} . Having made this replacement, we now view the exponent of $\log x$ in (13) as a function of β_{i-1} , thinking of α and ϵ as fixed. The minimum value of this function on the closed interval $[\beta, 1]$ occurs when $\beta_{i-1} = \alpha - \frac{1}{5}$, resulting in a contribution of

$$\ll_{\epsilon} x / (\log x)^{\frac{9}{5} + \alpha \log(\frac{1}{5}) - 3\epsilon}.$$

Since there are $O(\log_2 x)$ sub-cases, the contribution from all values of i is

$$\ll_{\epsilon} \frac{x}{(\log x)^{\frac{9}{5} + \alpha \log(\frac{1}{5}) - 4\epsilon}}. \tag{15}$$

Optimization. We now choose α, β to minimize the size of the total exceptional set obtained by adding the estimates (4), (6), (9), (15). (The other exceptional sets appearing in the argument are of total size $O(x/(\log x)^{1/2})$, which is tiny on the scale we are interested in, so we ignore these.) The optimal choice of α is obtained by setting the exponent $Q(\alpha)$ from (4) equal to the exponent $\frac{9}{5} + \alpha \log(\frac{1}{5})$ from (15), which yields $\alpha = 1.114478\dots$. This leads to the exponent $Q(\alpha) = 0.006316\dots$. Choosing $\beta = 0.7$, say, the remaining error terms (6) and (9) are negligible. (Note that (14) is satisfied for these choices of α and β , and that the various choices of the parameter z in the proof all satisfy $0 < z < 1$ as required.) Thus, our count is smaller than $x/(\log x)^{0.0063}$ for all sufficiently large values of x , which completes the proof of Theorem 1.

Remark. Our argument can be modified to show that the number of squarefull integers in $[1, x^2]$ which belong to \mathcal{V} is at most $x/(\log x)^{0.0063}$ once x is large. Indeed, all but $O(x/(\log x)^{1/2})$ squarefull numbers in $[1, x^2]$ are of the form $m^3 n^2$ with $m \leq \log x$. For each such m , we find that the number of n with $m^3 n^2 \in \mathcal{V} \cap [1, x^2]$ is $O(x \cdot m^{-3/2} / (\log x)^{0.006316})$, uniformly in m . Now we sum on m to get the claim.

4. A LOWER BOUND AND A HEURISTIC

4.1. Proof of Theorem 2.

Proof. A number $n \leq x$ has at most a bounded number of representations as $ap_1 p_2$ where p_1, p_2 is a pair of distinct primes at least $x^{1/10}$, in fact, at most 72 such (ordered) representations. Suppose y is a power of 2 with $x^{1/10} \leq y < x^{1/5}$. For each prime $p \in (y, 2y)$, let \mathcal{Q}_p denote the set of primes $q \leq x$ with $q \equiv 1 \pmod{p^2}$. From the Brun–Titchmarsh inequality, it follows that

$$\#\mathcal{Q}_p \ll \frac{x}{\varphi(p^2) \log(x/p^2)} \ll \frac{x}{y^2 \log x}. \tag{16}$$

This upper bound is usually correct as a lower bound as well, in fact it follows from [1, Theorem 2.1] that but for $O(1)$ choices for p (for all large x , uniformly in y),

$$\#\mathcal{Q}_p \geq \frac{x}{2\varphi(p^2) \log x} \gg \frac{x}{y^2 \log x}. \tag{17}$$

For a prime $p \in (y, 2y)$ and $q \in \mathcal{Q}_p$, we consider the integer $a = (q - 1)/p^2 < x/y^2$. In fact, for each integer $a < x/y^2$, let $N(a)$ denote the number of pairs of primes p, q with $p \in (y, 2y)$, $q \in \mathcal{Q}_p$, and $a = (q - 1)/p^2$. Thus, from (17),

$$\sum_{a < x/y^2} N(a) = \sum_{p \in (y, 2y)} \#\mathcal{Q}_p \gg \frac{y}{\log y} \cdot \frac{x}{y^2 \log x} \gg \frac{x}{y(\log x)^2},$$

and similarly, using (16), we have

$$\sum_{a < x/y^2} N(a) \ll \frac{x}{y(\log x)^2}.$$

It follows from Cauchy's inequality that

$$\sum_{a < x/y^2} N(a)^2 \geq \frac{y^2}{x} \left(\sum_{a < x/y^2} N(a) \right)^2 \gg \frac{y^2}{x} \cdot \frac{x^2}{y^2(\log x)^4} = \frac{x}{(\log x)^4}.$$

The last two displays imply that

$$\sum_{a < x/y^2} (N(a)^2 - N(a)) \gg \frac{x}{(\log x)^4}.$$

The expression $N(a)^2 - N(a)$ is the number of distinct pairs $(p_1, q_1), (p_2, q_2)$, where for $i = 1, 2$, p_i is a prime between y and $2y$, $q_i \in \mathcal{Q}_{p_i}$, and $(q_i - 1)/p_i^2 = a$. Since the pairs are distinct, it follows that $p_1 \neq p_2$ and $q_1 \neq q_2$. Also

$$\varphi(q_1 q_2) = (ap_1 p_2)^2 \quad \text{and} \quad ap_1 p_2 < a(\max\{p_1, p_2\})^2 < x.$$

Thus, each triple (a, p_1, p_2) generated by this process gives rise to a value $n = ap_1 p_2 < x$ for which $n^2 \in \mathcal{V}$. For a certain constant $c > 0$, there are at least $cx/(\log x)^4$ of these triples, uniformly in y . Summing y over the powers of 2 in $[x^{1/10}, x^{1/5})$ creates at least $c'x/(\log x)^3$ distinct triples of this kind. But the observation at the start of the proof shows that an individual $n = ap_1 p_2$ can arise from only $O(1)$ triples. Thus, $V_{\square}(x) \gg x/(\log x)^3$. \square

4.2. A heuristic. The above proof gives a lower estimate for the number of squares of the form $\varphi(q_1 q_2)$, where q_1, q_2 are distinct primes. One might ask what the “true” answer is, and more generally for the distribution of squares of the form $\varphi(m)$ where m is the product of k distinct odd primes, say $m = q_1 \cdots q_k$. Such a square n^2 has a natural factorization as $(q_1 - 1) \cdots (q_k - 1)$. If $q_i - 1$ is written as $a_i b_i^2$ with a_i squarefree, it follows that $a_1 \cdots a_k$ is a square. For the case $k = 2$, as we have seen in the proof above, this forces $a_1 = a_2$. In the case $k = 3$ we have three numbers A_1, A_2, A_3 with $a_i = A_1 A_2 A_3 / A_i$, for $i = 1, 2, 3$. The situation gets more complicated for 4 or more primes.

Suppose that a number $n \leq x$ is divisible by 4, $n/4$ is squarefree, and $\Omega(n/4) \geq \alpha \log_2 x$, where we fix a real number $\alpha > 1$. The number of ordered factorizations of n as $A_1 A_2 A_3 b_1 b_2 b_3$ with at least 2 of A_1, A_2, A_3 even is at least $6^{\Omega(n/4)} \geq (\log x)^{\alpha \log 6}$. The “chance” that each of $1 + b_i^2 A_1 A_2 A_3 / A_i$ is prime for $i = 1, 2, 3$ “should be” about $(\log x)^{-3}$. So, if $\alpha \log 6 > 3$, i.e., $\alpha > 3/\log 6$, there should be at least one such factorization. Thus, most numbers $n \leq x$ with $n/4$ squarefree and $\Omega(n/4) > \alpha \log_2 x$ with α a fixed real larger than $3/\log 6$ should have $n^2 \in \mathcal{V}$. It should then follow that $V_{\square}(x) \gg x/(\log x)^{Q(\alpha)}$. Since $Q(3/\log 6) = 0.18864255 \dots$, we thus should have $V_{\square}(x) \geq x/(\log x)^{0.189}$ for all sufficiently large values of x . Note that repeating this argument with products of 2 or 4 primes gives a worse result.

5. SQUARE VALUES OF THE SUM-OF-DIVISORS FUNCTION

Both Theorems 1 and 2 remain true with σ replacing φ . When porting over the proofs, the main idea is to replace every occurrence of $\varphi(q) = q - 1$ with $\sigma(q) = q + 1$. This works without much fuss for Theorem 2, and we leave the details to the reader. For Theorem 1, we meet additional difficulties owing to the more complicated behavior of σ on prime powers. In this section, we sketch a way around these roadblocks.

5.1. **Outline.** Assume that $n \leq x$ is such that $n^2 = \sigma(m)$. We can assume all of our previous conditions (i)–(v) on n and m , with the exception of (iv), which we replace with

(iv') m has no prime power divisor $q^e > \exp((\log x)^{1/2})$ with $e \geq 2$.

We leave the justification of (iv') to the end of this section, where it is shown (Lemma 9) that this assumption introduces an exceptional set of size $O(x/(\log x)^{1/4})$. For the rest of the argument, we fix the values of α and β to the constants we found above. Thus, $\alpha = 1.114478\dots$ and $\beta = 0.7$.

With $p = P(n)$, we have $p^2 \mid n^2 = \sigma(m)$. It cannot be the case that $p \mid \sigma(q^e)$ for a prime power $q^e \parallel m$ having $e \geq 2$, for then $2q^e > q^e + q^{e-1} + \dots + 1 = \sigma(q^e) \geq p$, forcing $q^e > \frac{p}{2} > \frac{1}{2}x^{1/2 \log_2 x}$ and contradicting (iv'). This leaves only two possibilities:

- I'. there are two different primes $q_1, q_2 \parallel m$ with $q_i \equiv -1 \pmod{p}$ for $i = 1, 2$,
- II'. there is a prime $q \parallel m$ with $q \equiv -1 \pmod{p^2}$.

Case I'. This is handled exactly as Case I above, replacing $q - 1$ with $q + 1$ throughout the argument. We find that the total count of n in Case I' satisfies our earlier upper bound (6).

Case II'. We start by writing $q + 1 = a(bp)^2$, so that $n = uabp$ for some integer u . Our first sub-case, when $P(ua) \leq \exp((\log x)^\beta)$, is handled exactly as was the first sub-case of Case II. Note that in the analogue of the sieve bound (7), the character χ_a appears in place of χ_{-a} . (We do not have to worry that a is a square, as that would imply $q = a(bp)^2 - 1$ factors.) This sub-case makes a total contribution of size (9).

In the remaining sub-cases, $P(ua) > \exp((\log x)^\beta)$. We again partition these according to the interval \mathcal{I}_i to which $p_2 := P(ua)$ belongs. Reasoning as in our treatment of Case II, we find that $p_2 \parallel n$; moreover, if we choose k so that $p_2^k \parallel q + 1$, then $k = 0$ or 1 according to whether or not $p_2 \mid a$. Hence,

$$p_2 \mid \frac{n^2}{q+1} = \sigma(m/q).$$

Thus, there is a prime power $q_2^e \parallel m/q$ for which p_2 divides $\sigma(q_2^e)$. Note that $q_2^e > \frac{1}{2}p_2 > \frac{1}{2} \exp((\log x)^\beta)$, so that if $e \geq 2$, we obtain a contradiction with (iv'). So $e = 1$ and $p_2 \mid q_2 + 1$. We choose j so that $p_2^j \parallel q_2 + 1$. Then $j = 1$ or $j = 2$, and $k + j \leq 2$. We now set $q_1 = q$, $p_1 = p$, $b_1 = b$, and continue to mimic our earlier arguments. We find that the contribution from all of the possible sub-cases of this sort satisfies (15).

Combining our estimates as before, we obtain the σ -analogue of Theorem 1 with the same exponent 0.0063.

5.2. Proof that we can assume (iv').

Lemma 9. *The count of $n \leq x$ with $n^2 = \sigma(m)$ for some m failing (iv') is $O(x/(\log x)^{1/4})$.*

Proof. We continue to assume that x is large. For the duration of the argument, we let $y = \exp((\log x)^{1/2})$. Suppose that $q^e \parallel m$. Then $\sigma(q^e) \mid \sigma(m) = n^2$, and so $r_{q^e} := \prod_{\ell^f \parallel \sigma(q^e)} \ell^{\lceil f/2 \rceil}$ is a divisor of n . Thus,

$$\frac{1}{x} \#\{n \leq x : n^2 = \sigma(m) \text{ for an } m \text{ where (iv') fails}\} \leq \sum^{(1)} + \sum^{(2)} + \sum^{(3)}, \quad (18)$$

where

$$\sum^{(1)} := \sum_{\substack{q^e > y \\ e \geq 3}} \frac{1}{r_{q^e}}, \quad \sum^{(2)} := \sum_{\substack{q > \sqrt{y} \\ r_{q^2} > q \log q}} \frac{1}{r_{q^2}}, \quad \text{and} \quad \sum^{(3)} := \sum_{\substack{q > \sqrt{y} \\ r_{q^2} \leq q \log q}} \frac{1}{r_{q^2}}.$$

Since $r_{q^e} \geq (\sigma(q^e))^{1/2} > q^{e/2}$, we have $\sum^{(1)} \leq \sum_{q^e > y, e \geq 3} q^{-e/2} \leq \sum_{\text{cubefull } c > y} c^{-1/2} \ll y^{-1/6}$, using in the final step that the count of cubefull numbers up to height t is $O(t^{1/3})$. By partial summation and the prime number theorem, $\sum^{(2)} \leq \sum_{q > \sqrt{y}} (q \log q)^{-1} \ll (\log y)^{-1}$. It remains to estimate $\sum^{(3)}$.

Let us show that $\mathcal{Q} := \{q : r_{q^2} \leq q \log q\}$ is a sparse set of primes. We begin with a simple observation: If $q^2 + q + 1$ has an exact prime divisor $\ell_0 > (\log q)^2$, then

$$r_{q^2} = \ell_0 \prod_{\substack{\ell^f \parallel q^2 + q + 1 \\ \ell \neq \ell_0}} \ell^{\lfloor f/2 \rfloor} \geq \ell_0 \sqrt{\frac{q^2 + q + 1}{\ell_0}} > q \sqrt{\ell_0} > q \log q,$$

and thus $q \notin \mathcal{Q}$. So if we suppose that $q \in \mathcal{Q} \cap (t/2, t]$ for a large real number t , then $q \in \mathcal{Q}_1 \cup \mathcal{Q}_2$, where

$$\mathcal{Q}_1 := \{q \in (t/2, t] : q^2 + q + 1 \text{ has no prime divisors in } ((\log t)^2, t^{1/10}]\},$$

$$\mathcal{Q}_2 := \{q \in (t/2, t] : \ell^2 \mid q^2 + q + 1 \text{ for some } \ell \in ((\log t)^2, t^{1/10}]\}.$$

Let $\varrho(r)$ be the number of roots modulo r of the polynomial $X^2 + X + 1$. For primes $\ell > 3$, we have $\varrho(\ell) = 2$ when $\ell \equiv 1 \pmod{3}$ and $\varrho(\ell) = 0$ otherwise. By the upper bound sieve (for instance, in the form of [14, Theorem 4.2, p. 134]),

$$\#\mathcal{Q}_1 \ll \frac{t}{\log t} \prod_{(\log t)^2 < \ell \leq t^{1/10}} \left(1 - \frac{\varrho(\ell)}{\ell}\right) \ll \frac{t}{(\log t)^2} \log_2 t \ll \frac{t}{(\log t)^{3/2}}.$$

(To estimate the product, we used a version of Mertens's theorem for primes congruent to 1 modulo 3.) We estimate $\#\mathcal{Q}_2$ crudely. Observing that $\varrho(\ell^2) \leq 2$ for all primes $\ell > 3$ (for instance, by Hensel's lemma), we obtain immediately that

$$\#\mathcal{Q}_2 \leq \sum_{(\log t)^2 < \ell \leq t^{1/10}} \left(\frac{2t}{\ell^2} + 2\right) \ll \frac{t}{(\log t)^2}.$$

Hence, $\#\mathcal{Q} \cap (t/2, t] \leq \#\mathcal{Q}_1 + \#\mathcal{Q}_2 \ll t/(\log t)^{3/2}$. Summing dyadically, we find that $\#\mathcal{Q} \cap [1, t] \ll t/(\log t)^{3/2}$ for all $t \geq 3$.

We now return to the problem of estimating $\sum^{(3)}$. Using the lower bound $r_{q^2} > q$, we find that $\sum^{(3)} \leq \sum_{q > \sqrt{y}, q \in \mathcal{Q}} q^{-1} \ll (\log y)^{-1/2}$, by partial summation. Lemma 9 now follows from (18) and our estimates for $\sum^{(1)}$, $\sum^{(2)}$, and $\sum^{(3)}$. \square

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