# ADDENDA: COMPUTING EUCLIDEAN BELYI MAPS 

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This note gives an addenda for the article Computing Euclidean Belyi maps [1].

## 1. Addenda

The addenda is summarized in the following additional remark.
Remark 3.2.10. If in Algorithm 2.4 .4 we compute instead the Smith normal form (SNF) of $A$ as $\left(\begin{array}{cc}n & 0 \\ 0 & m\end{array}\right)=P A Q$ (with $n \mid m$ ), the result gives a basis for $\Lambda_{\Gamma}$ relative to a basis for $\Lambda_{\Delta}$ such that $\Lambda_{\Gamma}=\left\langle n \omega_{1}^{\prime}, m \omega_{2}^{\prime}\right\rangle$ with $\Lambda_{\Delta}=\left\langle\omega_{1}^{\prime}, \omega_{2}^{\prime}\right\rangle$. Accordingly, we adjust Step 4 in Algorithm 3.2.5 by replacing the occurrences of $\omega_{1}$ and $\omega_{2}$ respectively with $\omega_{1}^{\prime}=a \omega_{1}+b \omega_{2}$ and $\omega_{2}^{\prime}=c \omega_{1}+d \omega_{2}$ where $Q^{-1}=\left(\begin{array}{ll}a & b \\ c & d\end{array}\right)$.

Incorporating Remark 3.2.9, we may further simplify by factoring $n$ from each entry in our basis matrix (corresponding to factoring the multiplication by $n$ map from $\widehat{\psi}$ ). This reduces us to the case $n=1$ in Algorithm 3.2.5.

In more detail, computing the map $\psi: E(\Gamma) \rightarrow E(\Delta)$ is the most complicated and costly step in Algorithm 3.5.1. To do so, we must first determine a basis for the lattice $\Lambda_{\Gamma}$ relative to a basis for the lattice $\Lambda_{\Delta}$. In Corollary 2.2.7, we make a "standard" choice for the basis vectors $\omega_{1}$ and $\omega_{2}$ for $\Lambda_{\Delta}$ that coincide with the periods Magma assigns to our canonical curves $E_{\square}$ and $E_{\square}$. Algorithm 2.4.4 then produces a two column matrix $A$ whose rows, taken as coordinates relative to the basis vectors $\omega_{1}$ and $\omega_{2}$, give a set of vectors that span $\Lambda_{\Gamma}$.

Reducing $A$ to Hermite normal form and taking its first two rows gives a basis matrix

$$
B_{H}:=\left(\begin{array}{cc}
n_{1} & n_{2} \\
0 & m_{2}
\end{array}\right)
$$

such that $\Lambda_{\Gamma}=\left\langle n_{1} \omega_{1}+n_{2} \omega_{2}, m_{2} \omega_{2}\right\rangle$.
If, instead, we reduce $A$ to Smith normal form and take its first two rows, we obtain a matrix of the form

$$
B_{S}:=\left(\begin{array}{cc}
n & 0 \\
0 & m
\end{array}\right)
$$

where $n$ divides $m$. Like with $B_{H}$, the matrix $B_{S}$ describes a basis for $\Lambda_{\Gamma}$ relative to a basis for $\Lambda_{\Delta}$ such that $\Lambda_{\Gamma}=\left\langle n \omega_{1}^{\prime}, m \omega_{2}^{\prime}\right\rangle$ with $\Lambda_{\Delta}=\left\langle\omega_{1}^{\prime}, \omega_{2}^{\prime}\right\rangle$. We note that the basis vectors $\omega_{1}^{\prime}$ and $\omega_{2}^{\prime}$ need not be the same as $\omega_{1}$ and $\omega_{2}$.

Because Magma's implementation of the Weierstrass $\wp$-function takes inputs relative to $\omega_{1}$ and $\omega_{2}$, it is then necessary to relate $\omega_{1}^{\prime}$ and $\omega_{2}^{\prime}$ back to the "standard" basis vectors. Let $P, Q \in \mathrm{GL}_{2}(\mathbb{Z})$ be such that $B_{S}=P B_{H} Q$. Then the matrices $P$ and $Q$ correspond, respectively, to elementary row and column operations performed on $B_{H}$ that transform it to $B_{S}$. As each elementary column operation
corresponds to an invertible change to the choice of basis for $\Lambda_{\Delta}$, we can recover the relationship between each $\omega_{i}^{\prime}$ and $\omega_{i}$ from the matrix $Q$ indicated above. Specifically, if

$$
Q^{-1}=\left(\begin{array}{ll}
a & b \\
c & d
\end{array}\right)
$$

then $\omega_{1}^{\prime}=a \omega_{1}+b \omega_{2}$ and $\omega_{2}^{\prime}=c \omega_{1}+d \omega_{2}$.
Working with $B_{S}$ rather than $B_{H}$ simplifies our computation of the isogeny $\psi: E(\Gamma) \rightarrow E(\Delta)$. Algorithm 3.2.5 describes a procedure for computing $\psi$ (by first computing its dual, $\widehat{\psi}$ ) that assumes the Hermite basis matrix $B_{H}$. If we instead work with the Smith matrix $B_{S}$, we may assume that $n_{2}=0$ and let $n_{1}=n$ and $m_{2}=m$. Incorporating remark 3.2.9 and recalling that $n$ divides $m$, we may further simplify by factoring $n$ from each entry in our basis matrix (corresponding to factoring the multiplication by $n$ map from $\widehat{\psi}$ ), leaving us with the matrix

$$
\frac{1}{n} B_{S}=\left(\begin{array}{cc}
1 & 0 \\
0 & m / n
\end{array}\right)
$$

where $m / n \in \mathbb{Z}$.
The combined effect of Remark 3.2.9 and this Smith simplification allows us to always assume in Algorithm 3.2.5 a basis matrix $B$ of particularly simple form:

$$
B:=\left(\begin{array}{ll}
1 & 0 \\
0 & d
\end{array}\right)
$$

This basis matrix gives coordinates relative to $\omega_{1}^{\prime}$ and $\omega_{2}^{\prime}$ rather than $\omega_{1}$ and $\omega_{2}$. Accordingly, we adjust the implementation of step 4 in Algorithm 3.2.5 by replacing the occurrences of $\omega_{1}$ and $\omega_{2}$ respectively with $\omega_{1}^{\prime}=a \omega_{1}+b \omega_{2}$ and $\omega_{2}^{\prime}=c \omega_{1}+d \omega_{2}$ as obtained above.

## References

[1] Matt Radosevich and John Voight, Computing Euclidean Belyi maps, accepted to J. Théorie Nombres Bordeaux.

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