## Math 38: Graph Theory

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7.1.26. Let G be a regular graph with cut-vertex. We will show that  $\chi'(G) > \Delta(G)$ .

Suppose that G is k-regular and that vertex  $v \in V(G)$  is a cut-vertex. Furthermore, assume that we have a proper k-edge-coloring of E(G) and thus  $\chi'(G) = k = \Delta(G)$ .

Vertex v has k edges incident with it. These edges are colored 1 through k. Suppose that the edge colored i has endpoints v and  $v_i$ . We will show that for all  $1 \le i < j \le k$ , vertices  $v_i$  and  $v_j$  are connected by a path in G - v.

Starting with vertex  $v_i$ , find a maximum walk whose edges are alternately colored i and j. This path must begin with color j, since  $v_i$  does not have an edge incident to it colored i in G - v. The walk can never return to a previously encountered vertex, otherwise we would have two edges colored the same that are incident to the same vertex, which would contradict the assumption that our edge-coloring is proper. And finally, each edge colored j leads us to a vertex that is incident with an edge colored i. This is due to the fact that we started at  $v_i$ , the only vertex that is not incident with an edge colored i, and we know we can never return to vertex  $v_i$ . Therefore, our walk ends when we reach a vertex that does not have an edge incident with it that is colored j. But  $v_j$  is the only such vertex. Thus our maximum walk must be a  $v_i, v_j$ -walk.

Therefore, any vertex in the component of G - v containing  $v_i$  is connected to any vertex in the component of G - v containing  $v_j$ . Since this is true for all i and j, we conclude that G - v has the same number of components as G. But this contradicts the assumption that v was a cut-vertex.

Therefore, if G is a k-regular graph with a cut-vertex then  $\chi'(G) > k = \Delta(G)$ , as promised.