

Barnett
4/18/13

SOLUTIONS

Math 56 Compu & Expt Math, Spring 2013: Midterm 1

4/18/13, pencil and paper, 2 hrs, 50 points. Good luck!

1. [9 points]

if you included $f'(x)$ error on input, messier but same outcome.

- [4] (a) Say a computer's algorithm for e^x has relative error in the output of up to ϵ_{mach} , for $-1 \leq x \leq 1$. Does this guarantee that the algorithm is backward stable in this domain?

We're told $\tilde{f}(x) = e^x (1 + \epsilon_1)$ for $|\epsilon_1| \leq \epsilon_{mach}$.

defn. $f(\tilde{x}) = e^{x(1+\epsilon)}$

for some $\epsilon = O(\epsilon_{mach})$

$= e^x e^{\epsilon x} \approx e^x (1 + \epsilon x + \dots)$
 $O(\epsilon^2)$, ignore

cancel e^x ; leaves $x + \epsilon_1 = x + \epsilon x$

so $\epsilon = \frac{\epsilon_1}{x}$ which is not $O(\epsilon_{mach})$ for $x \rightarrow 0$. \Rightarrow No.

- [3] (b) Repeat the question for $\sin x$ in the same domain.

Told $(\sin x)(1 + \epsilon_1)$ defn. blew stab $f(\tilde{x}) = \sin(x(1 + \epsilon))$

$\approx \sin x + \epsilon x \cos x + O(\epsilon^2)$
Taylor or addition thm.

subtract $\sin x$ from both sides:

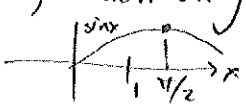
$\epsilon_1 \sin x = \epsilon x \cos x$

$\Rightarrow \epsilon = \frac{\tan x}{x} \epsilon_1$ so $\epsilon = O(\epsilon_{mach})$ uniformly over all $x \in [-1, 1]$
bounded in $-1 \leq x \leq 1$

\Rightarrow Yes,

- [2] (c) For some x outside $[-1, 1]$ one of the above algorithms cannot be backward stable. Which one, and for what x ?

We look for places where small changes in \tilde{x} fail to account for the known relative error in f . E.g. where $f'(x) = 0$, which only happens for $\sin x$ at $x = \frac{\pi}{2}(1 + 2n)$, eg. $x = \pi/2$

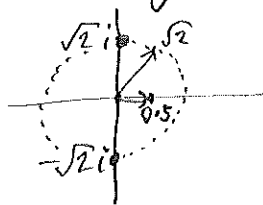


These are where $f'(x) = 0$.

2. [8 points] Consider $f(x) = 1/(2+x^2)$.

- [4] (a) What type, and order/rate, do you expect for convergence of the Taylor series truncated to terms less than x^n , expanding about the origin, when evaluated at $x = 0.5$? Explain

First way: find singularities in \mathbb{C} , ie $2+x^2=0$, $x_{\text{sing}} = \pm\sqrt{2}i$



Then on Taylor says exponentiated with rate $\frac{\text{center to } x \text{ dist}}{\text{center to } x_{\text{sing}} \text{ dist}}$

so rate $r = \frac{0.5}{\sqrt{2}} = \frac{1}{2\sqrt{2}}$

OR 2nd way: $f(x) = \frac{1}{2} \left(1 + \frac{x^2}{2}\right)^{-1} = \frac{1}{2} \left[1 - \frac{x^2}{2} + \frac{x^4}{2^2} - \dots \frac{x^n}{2^{n/2}} \dots \right]$ n even only.

tail of series $\epsilon_n = \sum_{\substack{k \geq n \\ k \text{ even}}} (\pm) \frac{x^k}{2^{k/2}}$, $|\epsilon_n| \leq \sum_{k \geq n} \frac{x^k}{2^{k/2}} = \left(\frac{x}{\sqrt{2}}\right)^n \sum_{k=0}^{\infty} \frac{x^k}{2^{k/2}} \leq C \left(\frac{x}{\sqrt{2}}\right)^n$

[1] Write an upper bound on the error reflecting this convergence, in big-O notation: $= C \frac{1}{(2\sqrt{2})^n}$

error up to x^n term, $\epsilon_n = O(r^n) = O\left(\frac{1}{(2\sqrt{2})^n}\right)$

- [3] (b) Estimate up to what power x^n is needed for this series to reach 16-digit accuracy.

Want $\epsilon_n \approx 10^{-16}$

ie $\frac{1}{(2\sqrt{2})^n} \approx 10^{-16}$

ie $n \ln(2\sqrt{2}) \approx \ln 10^{+16}$

$$n \approx \frac{\ln 10^{16}}{\ln(2\sqrt{2})} = \frac{\log_{10} 10^{16}}{\log_{10} 2\sqrt{2}} \approx \frac{16}{\frac{1}{2}} \approx 32$$

3. [8 points] Consider the "left-sided" finite-difference approximation $f'(x) \approx \frac{f(x) - f(x-h)}{h}$

- [5] (a) Derive a rigorous bound on the error that applies to each $h > 0$ [Hint: your bound will need to involve properties of f] *suggests Taylor's Theorem, so expand w/ remainder.*

$$\frac{f(x) - f(x-h)}{h} = \frac{f(x) - \left[f(x) - h f'(x) + \frac{h^2}{2} f''(q) \right]}{h}$$

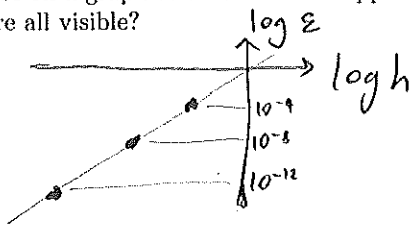
some $q \in [x-h, x]$

$$= \underbrace{f'(x)}_{\text{what we want}} - \underbrace{\frac{h}{2} f''(q)}_{\text{must be the error.}}$$

$$\Rightarrow \text{error } |\epsilon_h| \leq \frac{h}{2} \max_{q \in [x-h, x]} |f''(q)| = O(h)$$

note theorem required $f \in C^2([x-h, x])$

- [1] (b) What axes would one choose on a graph so that the error appears as a straight line and yet data at $h = 10^{-4}, 10^{-8}, 10^{-12}$ are all visible?



If you plot ϵ vs h it's also straight line, but if $h = 10^{-4}$ visible, others are crammed into the origin, invisible.

- [2] (c) Explain what happens to the error of the approximation in practice as $h \rightarrow 0$

In practice, f is computed to some relative error (maybe $O(\epsilon_{mach})$) so as h approaches $O(\epsilon_{mach})$ we get catastrophic cancellation and the approximation becomes terrible.

BONUS Roughly what h has the smallest error?

Balancing two sources of error, $O(h)$ from Taylor, & $O\left(\frac{\epsilon_{mach}}{h}\right)$ from evaluations of f , we get $h \approx \frac{\epsilon_{mach}}{h}$ so $h \approx \sqrt{\epsilon_{mach}} \approx 10^{-8}$, say.

matrix A

4. [7 points] Consider the linear system $\begin{bmatrix} 1 & 0 \\ 10^5 & 1 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} = \begin{bmatrix} b_1 \\ b_2 \end{bmatrix}$

How many digits accuracy (relative to the solution norm $\sqrt{x_1^2 + x_2^2}$) are you guaranteed in the solution if the system is solved by a backward stable algorithm with $\epsilon_{mach} = 10^{-16}$?

[You may assume a constant of 1 in the backward stability. Hint: full points for rigorous upper bound on the error; generous partial credit for intelligent estimates or other bounds]

We need $\kappa(A) = \|A\| \cdot \|A^{-1}\|$

ie, lower tells us to look for upper bound on $\kappa(A)$, but example pairs of vectors x, Ax only can give lower bounds. \Rightarrow Need exact κ calc.

$A^T A = \begin{bmatrix} 1 & 10^5 \\ 0 & 1 \end{bmatrix} \begin{bmatrix} 1 & 0 \\ 10^5 & 1 \end{bmatrix} = \begin{bmatrix} 10^{10} + 1 & 10^5 \\ 10^5 & 1 \end{bmatrix}$

\hookrightarrow eivals $\lambda^2 - (10^{10} + 2)\lambda - 10^{10} - 1 + (10^5)^2 = 0$ so $\lambda = \frac{1}{2}(10^{10} + 2 \pm \sqrt{(10^{10} + 2)^2 - 4})$

use $\sqrt{(10^{10} + 2)^2 - 4} \leq 10^{10} + 2$

so $\lambda_{max} \leq 10^{10} + 2$

so $\|A\| = \sqrt{\lambda_{max}(A^T A)} \leq \sqrt{10^{10} + 2}$

$A^{-1} = \frac{1}{ad-bc} \begin{bmatrix} d & -b \\ -c & a \end{bmatrix} = \begin{bmatrix} 1 & -10^5 \\ 10^5 & 1 \end{bmatrix}$

which gives identical calc as before $\Rightarrow \|A^{-1}\| \leq \sqrt{10^{10} + 2}$

So $\kappa(A) \leq 10^{10} + 2$

Backward Stability Theorem: rel. err $\leq \kappa \cdot O(\epsilon_{mach})$

so rel. err $\leq 10^{-6}$

$\approx 10^{10}$ \uparrow 10^{-16} if const=1

6-digit accuracy

Hard! Makes you think.

[BONUS] Find a right-hand side b for which the above worst-case prediction is (nearly) achieved.

$\kappa(A)$ is a worst-case value of $\kappa(b)$ for the linear solve problem

$\vec{x} = f(b) = A^{-1}b$, over all b . We need to rederive κ :

$\kappa(b) = \max_{\delta b \neq 0} \frac{\|\delta f\| / \|f\|}{\|\delta b\| / \|b\|} = \max_{\delta b \neq 0} \frac{\|A^{-1}\delta b\|}{\|\delta b\|} \cdot \frac{\|b\|}{\|A^{-1}b\|}$

$=: \|A^{-1}\|$, fixed

\rightarrow RHS-dependent part: want to maximize $\Rightarrow A^{-1}$ must shrink b the most, ie A must grow \vec{x} the most.

Easy to find \vec{x} which close to max growth:

$\begin{bmatrix} 1 & 0 \\ 10^5 & 1 \end{bmatrix} \begin{bmatrix} 1 \\ 0 \end{bmatrix} = \begin{bmatrix} 1 \\ 10^5 \end{bmatrix}$

\vec{x} b

so, this RHS does it.

Note, strangely, a 'typical' b will give $\vec{x} = A^{-1}b$ about 10^5 x larger, (not smaller), so rel. err $\approx O(\epsilon_{mach})!$

5. [7 points] Given $y > 0$, you wish to approximate $x = \sqrt{y}$ using elementary operations.

[1] (a) Derive a Newton iteration that converges to the desired x [Hint: x must be a root of something]

As in class pick $f(x) = x^2 - y$

$\overbrace{f(x)}$
not $x - \sqrt{y}$ which
demands computing \sqrt{y} !

$$\begin{aligned} x_{n+1} &= x_n - \frac{f(x_n)}{f'(x_n)} = x_n - \frac{x_n^2 - y}{2x_n} \\ &= \frac{x_n}{2} + \frac{y}{2x_n} \end{aligned}$$

[3] (b) Derive a big-O estimate on the error ε_n after n iterations.

I also gave 2/3 for
instead proving this.

This is hard, to make you think.

Recall Newton has 'quadratic' convergence, i.e. $\varepsilon_{n+1} \leq C \varepsilon_n^2$

$$\begin{aligned} \text{Start at } \varepsilon_0 : \quad \varepsilon_1 &\leq C \varepsilon_0^2 \\ \varepsilon_2 &\leq C \varepsilon_1^2 \leq C (C \varepsilon_0^2)^2 = C^3 \varepsilon_0^4 \\ \varepsilon_3 &\leq C (C (C \varepsilon_0^2)^2)^2 = C^7 \varepsilon_0^8 \end{aligned}$$

See by induction that etc. $\varepsilon_n \leq C^{2^n - 1} \varepsilon_0^{2^n} = \frac{1}{C} (C \varepsilon_0)^{2^n}$

so $\varepsilon_n = O(C^{2^n})$ for some const $C < 1$

↑
we hope!
(relies on good
starting guess,
 ε_0 small)

6. [11 points] Short answers.

[3]

(a) Prove whether $N + \frac{N}{(\log_{10} N) - 7} = O(N)$, giving, if true, a constant and corresponding condition on N .

$$\frac{f(N)}{g(N)} = 1 + \frac{1}{\log_{10} N - 7} \stackrel{?}{\leq} C$$

only bounded once $\log_{10} N > 7$

we assume $N \rightarrow \infty$
(forgot to say this).
holds $\forall N \geq 10^8$
with $C = 2$

[2]

(b) Prove whether $\sqrt{1+x^2} \sin x = o(x)$ as $x \rightarrow \infty$

$$\lim_{x \rightarrow \infty} \frac{f(x)}{g(x)} = \lim_{x \rightarrow \infty} \frac{\sqrt{1+x^2}}{x} \cdot \sin x$$

has no limit (despite being bounded)!

has limit = 1

[2]

(c) How close to 1 does x have to be such that the relative condition number of computing $\sqrt{x-1}$ is 10^8 ?

$$\kappa(x) = \left| \frac{x f'(x)}{f(x)} \right| = \left| \frac{x (x-1)^{-1/2} / 2}{\sqrt{x-1}} \right| = \left| \frac{x}{2(x-1)} \right|$$

Set $\kappa \approx 10^8$:

$$|x| = 2 \times 10^8 |x-1| \quad \text{so } |x-1| \approx \frac{1}{2 \times 10^8} \text{ ie } 5 \times 10^{-9}$$

distance away.

[2]

(d) Prove that $\|A^{-1}\| = \left(\min_{x \neq 0} \frac{\|Ax\|}{\|x\|} \right)^{-1}$ now pick $\tilde{x} = A^{-1}y$.

$$\|A^{-1}\| := \max_{y \neq 0} \frac{\|A^{-1}y\|}{\|y\|} = \max_{\tilde{x} \neq 0} \frac{\|\tilde{x}\|}{\|A\tilde{x}\|} = \left(\min_{\tilde{x} \neq 0} \frac{\|A\tilde{x}\|}{\|\tilde{x}\|} \right)^{-1}$$

since min of the reciprocal is reciprocal of the max.

[2]

(e) Give a definition of an algorithm $\tilde{f}(x)$ for a problem $f(x)$ being stable:

For each x in some domain of interest, there exists $\tilde{x} = x(1+\varepsilon)$

for some $\varepsilon = O(\varepsilon_{\text{mach}})$, such that $\frac{\|\tilde{f}(x) - f(\tilde{x})\|}{\|f(\tilde{x})\|} = O(\varepsilon_{\text{mach}})$.

"nearly right answer to nearly right question"

Or, for each x , $\exists \varepsilon = O(\varepsilon_{\text{mach}})$ st. $\|\tilde{f}(x) - f(x(1+\varepsilon))\| \leq \|f(x(1+\varepsilon))\| O(\varepsilon_{\text{mach}})$