## Math 38. Graph Theory.

## Solutions to Homework 1.

Sec 1.1, \# 10. It is true. Let $G$ be a simple disconnected graph. Take $u, v \in V(G)$ so that there is no $u, v$-path in $G$. In particular, we have $u v \notin E(G)$, and so $u v \in E(\bar{G})$.
To show that $\bar{G}$ is connected, it is enough to show that for every vertex $w$, there is a path from $u$ to $w$ in $\bar{G}$. To see this, take an arbitrary vertex $w$. If $u w \in E(\bar{G})$ we are done. Otherwise, the fact that $u w \in E(G)$ and that there is no $u, v$-path in $G$ forces $v w \notin E(G)$, that is, $v w \in E(\bar{G})$. But now the edges $u v$ and $v w$ give a path from $u$ to $w$ in $\bar{G}$.

Sec 1.1, \# 13. Yes. This is solved in example 1.3.8.
Sec 1.1, \# 16. No. Their complements are $C_{4}+C_{4}$ (a disjoint union of two 4 -cyles) and $C_{8}$, respectively, which are not isomorphic.

Sec 1.1, \# 29. Consider mutual acquaintances as edges of a 6 -vertex graph $G$ (and so mutual strangers are edges of $\bar{G})$. The problem asks to show that $G$ has either a clique of size 3 or an independent set of size 3 .

Take any $v \in G$. Among the remaining 5 vertices, either we have at least 3 adjacent to $v$ or at least 3 non-adjacent to $v$, by the pigeonhole principle. We can assume without loss of generality that at least 3 are adjacent to $v$ (if we are in the other case, we argue similarly using $\bar{G}$ instead of $G$ ), that is, $v u_{1}, v u_{2}, v u_{3} \in E(G)$. If one of $u_{1} u_{2}, u_{2} u_{3}, u_{1} u_{3}$ is in $E(G)$, then $G$ has a triangle, consisting of the two endpoints of this edge together with $v$. If not, then $G$ has an independent set $\left\{u_{1}, u_{2}, u_{3}\right\}$.

Sec 1.1, \# 30. Let $a_{i}$ denote the $i$-th column of the adjacency matrix $A$. Then the entry in row $i$ and column $i$ of $A^{2}$ equals $a_{i}^{T} a_{i}$, which is the number of ones in $a_{i}$, which is the degree of the $i$-th vertex $v_{i}$.
Similarly, letting $b_{i}^{T}$ denote the $i$-th row of the incidence matrix $M$, the entry in row $i$ and column $i$ of $M M^{T}$ equals $b_{i}^{T} b_{i}$, which is the number of ones in $b_{i}$, which is again degree of the $v_{i}$.
By a similar argument, the entry in position $i, j$ of $A^{2}$ is the number of common neighbors of $v_{i}$ and $v_{j}$, which equals the number of walks of length 2 from $v_{i}$ to $v_{j}$.
The entry in position $i, j$ of $M M^{T}$ is the number of edges between $v_{i}$ and $v_{j}$, which must be 0 or 1 for a simple graph.

