Theorem

Let G be a graph with n vertices. Then the following are equivalent:

- A) G is a tree (i.e., it is connected and has no cycles).
- B) G is connected and has n 1 edges.
- C) G has n 1 edges and no cycles.
- D) For any $u, v \in V(G)$, there is exactly one u, v-path.

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Some consequences:

- Every edge of a tree is a cut-edge.
- Adding an edge to a tree forms exactly one cycle.

A spanning tree of G is a subgraph of G with vertex set V(G) that is a tree.

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Proposition

If T and T' are spanning trees of G and $e \in E(T) \setminus E(T')$, then there is an edge $e' \in E(T') \setminus E(T)$ so that T - e + e' is a spanning tree of G.

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$$\operatorname{rad}(G) = \min_{u \in V(G)} \epsilon(u).$$

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Definition

The **center** of *G* is the subgraph induced by the vertices of minimum eccentricity.