## Summary of notation

$\alpha(G)=$ maximum size of an independent set in $G$
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\begin{gathered}
\beta^{\prime}(G) \geq \alpha(G) \\
\beta^{\prime}(G) \geq \frac{n(G)}{2} \geq \alpha^{\prime}(G)
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## More relations

## Theorem

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## Corollary

If $G$ is bipartite with no isolated vertices, then

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\alpha(G)=\beta^{\prime}(G) .
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Chapter 4
Connectivity and paths

## Connectivity

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In this chapter, graphs will have no loops (since loops do not affect connectivity).

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