

Cut-edges and cut-vertices

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- For $S \subseteq V(G)$, $G - S$ is the graph obtained by deleting the vertices in the set S from G .

Induced subgraphs

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Letting $\bar{T} = V(G) \setminus T$, we have that $G[T] = G - \bar{T}$.

Characterization of cut-edges

Theorem

An edge is a cut-edge if and only if it belongs to no cycle.

Characterization of bipartite graphs

We will characterize bipartite graphs in terms of their cycles.

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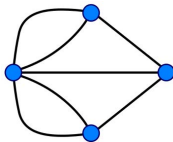
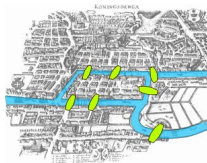
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Theorem (König 1936)

A graph is bipartite if and only if it contains no odd cycle.

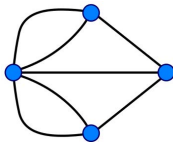
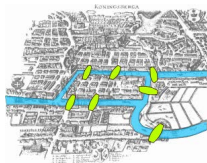
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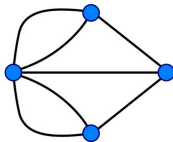
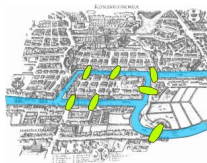


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A graph is **Eulerian** if it has a closed walk that contains each edge exactly once. Such a walk is called an **Eulerian circuit**.

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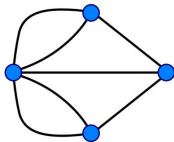
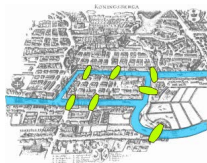
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If every vertex of a graph has degree ≥ 2 , then it contains a cycle.

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Recall the Königsberg bridge problem.



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Lemma

If every vertex of a graph has degree ≥ 2 , then it contains a cycle.

Theorem (Characterization of Eulerian graphs)

A connected graph is Eulerian if and only if all its vertices have even degree.