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Theorem (Characterization of Eulerian graphs)

A connected graph is Eulerian if and only if all its vertices have even degree.

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Definition

The **neighborhood** of v, denoted by $N_G(v)$ or N(v), is the set of vertices that are adjacent to v.

Proposition (Degree-sum formula)

$$\sum_{\in V(G)} d(v) =$$

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- A k-regular graph with n vertices has $\frac{nk}{2}$ edges.

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- Every graph has an even number of odd-degree vertices.
- A k-regular graph with n vertices has nk
 n particular, if n is odd, then k must be even.

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Proposition

A k-regular bipartite graph has the same number of vertices in each partite set.

Conjecture (The Reconstruction Conjecture)

If G is a simple graph with at least 3 vertices, then G is uniquely determined by the list of unlabeled subgraphs obtained from G by deleting one vertex.

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[Example]