

Chapter 2

Trees and distance

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Lemma

Every tree has at least two leaves.

Deleting a leaf from a tree produces another tree.

Theorem

Let G be a graph with n vertices. Then the following are equivalent:

- A) G is a tree (i.e., it is connected and has no cycles).*
- B) G is connected and has $n - 1$ edges.*
- C) G has $n - 1$ edges and no cycles.*
- D) For any $u, v \in V(G)$, there is exactly one u, v -path.*

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Some consequences:

- Every edge of a tree is a cut-edge.
- Adding an edge to a tree forms exactly one cycle.